

# TWO-SIDED ASSEMBLY LINE

## *Estimation of Final Results*

Waldemar Grzechca

*Institute of Automatic Control, The Silesian University of Technology, ul.Akademicka 16, 44-100 Gliwice, Poland*

**Keywords:** Assembly line balancing, Two-sided structure, Line time, Line efficiency, Smoothness index.

**Abstract:** The paper considers simple assembly line balancing problem and two-sided assembly line structure. In the last four decades a large variety of heuristic and exact solutions procedures have been proposed to balance one-sided assembly line in the literature. Some heuristic were given to balance two-sided lines, too. Some measures of solution quality have appeared in line balancing literature: balance delay (BD), line efficiency (LE), line time (LT) and smoothness index (SI). These measures are very important for estimation the balance solution quality. Author of this paper modified and discussed the line time and smoothness for two-sided assembly line. Some problems, which appeared during evaluations, are mentioned.

## 1 INTRODUCTION

The manufacturing assembly line was first introduced by Henry Ford in the early 1900's. It was designed to be an efficient, highly productive way of manufacturing a particular product. The basic assembly line consists of a set of workstations arranged in a linear fashion, with each station connected by a material handling device. The basic movement of material through an assembly line begins with a part being fed into the first station at a predetermined feed rate. A station is considered any point on the assembly line in which a task is performed on the part. These tasks can be performed by machinery, robots, and/or human operators. Once the part enters a station, a task is then performed on the part, and the part is fed to the next operation. The time it takes to complete a task at each operation is known as the process time (Sury, 1971). The cycle time of an assembly line is predetermined by a desired production rate. This production rate is set so that the desired amount of end product is produced within a certain time period (Baybars, 1986). For instance, the production rate might be set at 480 parts per day. Assuming an eight-hour shift, this translates into a requirement of 60 parts per hour (1 part per minute) being produced by the assembly line. In order for the assembly line to maintain a certain production rate, the sum of the processing times at each station must not exceed the stations' cycle time (Fonseca et. al, 2005). If the sum of the processing times within a station is less than the

cycle time, idle time is said to be present at that station (Erel, Erdal and Sarin, 1998). One of the main issues concerning the development of an assembly line is how to arrange the tasks to be performed. This arrangement may be somewhat subjective, but has to be dictated by implied rules set forth by the production sequence (Kao, 1976). For the manufacturing of any item, there are some sequences of tasks that must be followed. The assembly line balancing problem (ALBP) originated with the invention of the assembly line. Helgeson et. al (Helgeson and Birnie, 1961) were the first to propose the ALBP, and Salveson (Salveson, 1955) was the first to publish the problem in its mathematical form. However, during the first forty years of the assembly line's existence, only trial-and-error methods were used to balance the lines (Erel, Erdal and Sarin, 1998). Since then, there have been numerous methods developed to solve the different forms of the ALBP. Salveson (Salveson, 1955) provided the first mathematical attempt by solving the problem as a linear program. Gutjahr and Nemhauser (Gutjahr and Nemhauser, 1964) showed that the ALBP problem falls into the class of NP-hard combinatorial optimization problems. This means that an optimal solution is not guaranteed for problems of significant size. Therefore, heuristic methods have become the most popular techniques for solving the problem.

## 2 TWO-SIDED ASSEMBLY LINE

Two-sided assembly lines are typically found in producing large-sized products, such as trucks and buses. Assembling these products is in some respects different from assembling small products. Some assembly operations prefer to be performed at one of the two sides (Bartholdi, 1993).

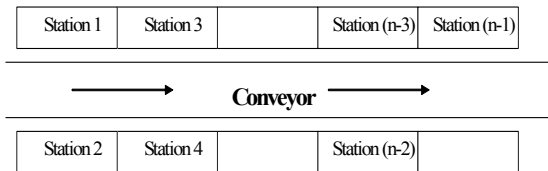


Figure 1: Two-sided assembly line.

Let us consider, for example, a truck assembly line. Installing a gas tank, air filter, and toolbox can be more easily achieved at the left-hand side of the line, whereas mounting a battery, air tank, and muffler prefers the right-hand side. Assembling an axle, propeller shaft, and radiator does not have any preference in their operation directions so that they can be done at any side of the line. The consideration of the preferred operation directions is important since it can greatly influence the productivity of the line, in particular when assigning tasks, laying out facilities, and placing tools and fixtures in a two-sided assembly line (Kim et. al, 2001). A two-sided assembly line in practice can provide several substantial advantages over a one-sided assembly line (Bartholdi, 1993). These include the following: (1) it can shorten the line length, which means that fewer workers are required, (2) it thus can reduce the amount of throughput time, (3) it can also benefit from lowered cost of tools and fixtures since they can be shared by both sides of a mated-station, and (4) it can reduce material handling, workers movement and set-up time, which otherwise may not be easily eliminated. These advantages give a good reason for utilizing two-sided lines for assembling large-sized products.

A line balancing problem is usually represented by a precedence diagram as illustrated in Figure 2. A circle indicates a task, and an arc linking two tasks represents the precedence relation between the tasks. Each task is associated with a label of  $(t_i, d)$ , where  $t_i$  is the task processing time and  $d$  (=L, R or E) is the preferred operation direction. L and R, respectively, indicate that the task should be assigned to a left- and a right-side station. A task associated with E can be performed at either side of the line.

While balancing assembly lines, it is generally needed to take account of the features specific to the lines. In a one-sided assembly line, if precedence relations are considered appropriately, all the tasks assigned to a station can be carried out continuously without any interruption. However, in a two-sided assembly line, some tasks assigned to a station can be delayed by the tasks assigned to its companion (Bartholdi, 1993). In other words, idle time is sometimes unavoidable even between tasks assigned to the same station. Consider, for example, task  $j$  and its immediate predecessor  $i$ . Suppose that  $j$  is assigned to a station and  $i$  to its companion station. Task  $j$  cannot be started until task  $i$  is completed. Therefore, balancing such a two-sided assembly line, unlike a one-sided assembly line, needs to consider the sequence-dependent finish time of tasks.

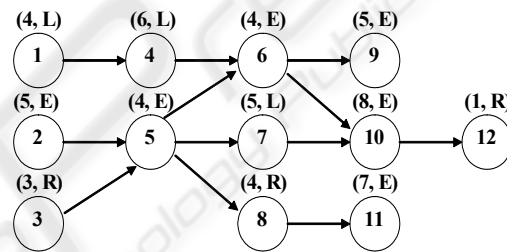


Figure 2: Precedence graph (cycle time =16).

This notion of sequence dependency further influences the treatment of cycle time constraint. Every task assigned to a station must be able to be completed within a predetermined cycle time. In a one-sided assembly line, this can readily be achieved by checking the total operation time of tasks assigned to a station. Therefore, a task not violating any precedence constraints can be simply added to the station if the resulting total amount of operation time does not exceed the cycle time. However, in a two-sided assembly line, due to the above sequence-dependent delay of tasks, the cycle time constraint should be more carefully examined. The amount of time required to perform tasks allocated to a station is determined by the task sequences in both sides of the mated-station as well as their operation time. It should be mentioned that two-sided assembly line is a special case of single assembly line. Therefore it is possible to use some procedures and measurements, which were for simple assembly line developed.

### 3 HEURISTIC APPROACH

#### 3.1 Grouping Tasks

A task group consists of a considered task  $i$  and all of its predecessors. Such groups are generated for every un–assigned task. As mentioned earlier, balancing a two–sided assembly line needs to additionally consider operation directions and sequence dependency of tasks, while creating new groups (Kim et. al, 2005).

While forming initial groups  $IG(i)$ , the operation direction is being checked all the time. It's disallowed for a group to contain tasks with preferred operation direction from opposite sides. But, if each task in initial group is E – task, the group can be allocated to any side. In order to determine the operation directions for such groups, the rules (direction rules DR) are applied:

DR 1. Set the operation direction to the side where tasks can be started earlier.

DR 2. The start time at both sides is the same, set the operation direction to the side where it's expected to carry out a less amount of tasks (total operation time of unassigned L or R tasks).

Generally, tasks resulting from “repeatability test” are treated as starting ones. But there is exception in form of first iteration, where procedure starts from searching tasks (initial tasks IT), which are the first ones in precedence relation. After the first step in the first iteration we get:

$$IG(1) = \{1\}, \text{Time}\{IG(1)\} = 2, \text{Side}\{IG(1)\} = 'L'$$

$$IG(2) = \{2\}, \text{Time}\{IG(2)\} = 5, \text{Side}\{IG(2)\} = 'E'$$

$$IG(3) = \{3\}, \text{Time}\{IG(3)\} = 3, \text{Side}\{IG(3)\} = 'R'$$

where:

$\text{Time}\{IG(i)\}$  – total processing time of  $i^{\text{th}}$  initial group,

$\text{Side}\{IG(i)\}$  – preference side of  $i^{\text{th}}$  initial group.

To those who are considered to be the first, the next tasks will be added, (these ones which fulfil precedence constraints).

Whenever new tasks are inserted to the group  $i$ , the direction, cycle time and number of immediate predecessors are checked. If there are more predecessors than one, the creation of initial group  $j$  comes to the end.

First iteration – second step

$$IG(1) = \{1, 4, 6\}, \text{Time}\{IG(1)\} = 8, \text{Side}\{IG(1)\} = 'L'$$

$$IG(2) = \{2, 5\}, \text{Time}\{IG(2)\} = 9, \text{Side}\{IG(2)\} = 'E'$$

$$IG(3) = \{3, 5\}, \text{Time}\{IG(3)\} = 7, \text{Side}\{IG(3)\} = 'R'$$

When set of initial groups is created, the last elements from those groups are tested for repeatability. If last element in set of initial groups  $IG$  will occur more than once (groups pointed by arrows), the groups are intended to be joined – if total processing time (summary time of considered groups) is less or equal to cycle time. Otherwise, these elements are deleted.

In case of occurring only once, the last member is being checked if its predecessors are not contained in Final set FS. If not, it's removed as well. So far, FS is empty.

First iteration – third step

$$IG(1) = \{1, 4\}, \text{Time}\{IG(1)\} = 4, \text{Side}\{IG(1)\} = 'L'$$

$$IG(2) = \{2, 3, 5\}, \text{Time}\{IG(2)\} = 12, \text{Side}\{IG(2)\} = 'R'$$

Whenever two or more initial groups are joined together, or when initial group is connected with those one coming from Final set – the “double task” is added to initial tasks needed for the next iteration. In the end of each iteration, created initial groups are copied to FS.

First iteration – fourth step

$$FS = \{ (1, 4); (2, 3, 5) \},$$

$$\text{Side}\{FS(1)\} = 'L', \text{Side}\{FS(2)\} = 'R'$$

$$\text{Time}\{FS(2)\} = 12, \text{Time}\{FS(1)\} = 14,$$

$$IT = \{5\}.$$

In the second iteration, second step, we may notice that predecessor of last task coming from  $IG(1)$  is included in Final Set,  $FS(2)$ . The situation results in connecting both groups under holding additional conditions:

$$\text{Side}\{IG(1)\} = \text{Side}\{FS(2)\},$$

$$\text{Time} + \text{time} < \text{cycle}.$$

After all, there is no more IT tasks, hence, preliminary process of creating final set is terminated.

The presented method for finding task groups is to be summarized in simplified algorithm form. Let  $U$  denote to be the set of un – assigned tasks yet and

$IG_i$  be a task group consisting of task  $i$  and all its predecessors (excluded from  $U$  set).

STEP 1. If  $U = \text{empty}$ , go to step 5, otherwise, assign starting task from  $U$ .

STEP 2. Identify  $IG_i$ . Check if it contains tasks with both left and right preference operation direction, then remove task  $i$ .

STEP 3. Assign operation direction  $Side\{IG_i\}$  of group  $IG_i$ . If  $IG_i$  has R-task (L-task), set the operation direction to right (left). Otherwise, apply so called direction rules DR.

STEP 4. If the last task  $i$  in  $IG_i$  is completed within cycle time, the  $IG_i$  is added to Final set of candidates  $FS(i)$ . Otherwise, exclude task  $i$  from  $IG_i$  and go to step 1.

STEP 5. For every task group in  $FS(i)$ , remove it from  $FS$  if it is contained within another task group of  $FS$ .

The resulting task groups become candidates for the mated-station.

$FS = \{(1,4), (2,3,5,8)\}$ .

### 3.2 Groups Assignment

The candidates are produced by procedures presented in the previous section, which claim to not violate precedence, operation direction restrictions, and what's more it exerts on groups to be completed within preliminary determined cycle time. Though, all of candidates may be assigned equally, the only one group may be chosen. Which group it will be – for this purpose the rules helpful in making decision, will be defined and explained below:

AR 1. Choose the task group  $FS(i)$  that may start at the earliest time.

AR 2. Choose the task group  $FS(i)$  that involves the minimum delay.

AR 3. Choose the task group  $FS(i)$  that has the maximum processing time.

In theory, for better understanding, we will consider a left and right side of mated – station, with some tasks already allocated to both sides. In order to achieve well balanced station, the AR 1 is applied, cause the unbalanced station is stated as the one which would probably involve more delay in future assignment. This is the reason, why minimization number of stations is not the only goal, there are also indirect ones, such as reduction of unavoidable delay. This rule gives higher priority to the station, where less tasks are allocated. If ties occurs, the AR 2 is executed, which chooses the group with the least amount of delay among the considered ones. This rule may also result in tie. The last one, points at

relating work within individual station group by choosing group of task with highest processing time.

For the third rule the tie situation is impossible to obtain, because of random selection of tasks. The implementation of above rules is strict and easy except the second one. Shortly speaking, second rule is based on the test, which checks each task consecutively, coming from candidates group  $FS(i)$  – in order to see if one of its predecessors have already been allocated to station. If it has, the difference between starting time of considered task and finished time of its predecessor allocated to companion station is calculated. The result should be positive, otherwise time delay occurs.

### 3.3 Final Procedures

Having rules for initial grouping and assigning tasks described in previous sections, we may proceed to formulate formal procedure of solving two – sided assembly line balancing problem (Kim et. al, 2005).

Let us denote companion stations as  $j$  and  $j'$ ,

$D(i)$  – the amount of delay,

$Time(i)$  – total processing time ( $Time\{FS(i)\}$ ),

$S(j)$  – start time at station  $j$ ,

STEP 1. Set up  $j = 1, j' = j + 1, S(j) = S(j') = 0, U$  – the set of tasks to be assigned.

STEP 2. Start procedure of group creating (3.2), which identifies

$FS = \{FS(1), FS(2), \dots, FS(n)\}$ . If  $FS = \emptyset$ , go to step 6.

STEP 3. For every  $FS(i), i = 1, 2, \dots, n$  – compute  $D(i)$  and  $Time(i)$ .

STEP 4. Identify one task group  $FS(i)$ , using AR rules in Section 3.3

STEP 5. Assign  $FS(i)$  to a station  $j$  ( $j'$ ) according to its operation direction, and update  $S(j) = S(j) + Time(i) + D(i)$ .  $U = U - \{FS(i)\}$ , and go to STEP 2.

STEP 6. If  $U \neq \emptyset$ , set  $j = j' + 1, j' = j + 1, S(j) = S(j') = 0$ , and go to STEP 2, Otherwise, stop the procedure.

## 4 MEASURES OF FINAL RESULTS OF ASSEMBLY LINE BALANCING PROBLEM

Some measures of solution quality have appeared in line balancing problem. Below are presented three of them (Scholl, 1998).

Line efficiency (LE) shows the percentage utilization of the line. It is expressed as ratio of total station time to the cycle time multiplied by the number of workstations:

$$LE = \frac{\sum_{i=1}^K ST_i}{c \cdot K} \cdot 100\% \quad (1)$$

where:

K - total number of workstations,  
c - cycle time.

Smoothness index (SI) describes relative smoothness for a given assembly line balance. Perfect balance is indicated by smoothness index 0. This index is calculated in the following manner:

$$SI = \sqrt{\sum_{i=1}^K (ST_{max} - ST_i)^2} \quad (2)$$

where:

ST<sub>max</sub> = maximum station time (in most cases cycle time),  
ST<sub>i</sub> = station time of station i.

Time of the line (LT) describes the period of time which is need for the product to be completed on an assembly line:

$$LT = c \cdot (K - 1) + T_K \quad (3)$$

where:

c - cycle time,  
K -total number of workstations.

### 5 NUMERICAL RESULTS

The results of proposed procedure for the example from Figure 2 are given in a Gantt chart – Figure 3. Before presenting performance measures for current example, it would be like to stress difference in estimation of line time form, resulting from restrictions of parallel stations. In two – sided line method within one mated-station, tasks are intended to perform its operations at the same time, as it is shown in example in Figure.3, where tasks 7, 11 respectively are processed simultaneously on single station 3 and 4, in contrary to one – sided heuristic methods. Hence, modification has to be introduced to that particular parameter which is the consequence of parallelism. Having two mated-stations from Figure 3, the line time LT is not 3\*16 + 13, as it was in original expression. We must treat those stations as two double ones (mated-stations), rather than individual ones S<sub>k</sub>. Accepting this line of reasoning, new formula is presented below:

$$LT = c \cdot (Km - 1) + \text{Max}\{t(S_K), t(S_{K-1})\} \quad (4)$$

where:

Km – number of mated-stations  
K – number of assigned single stations  
t(S<sub>k</sub>) – processing time of the last single station

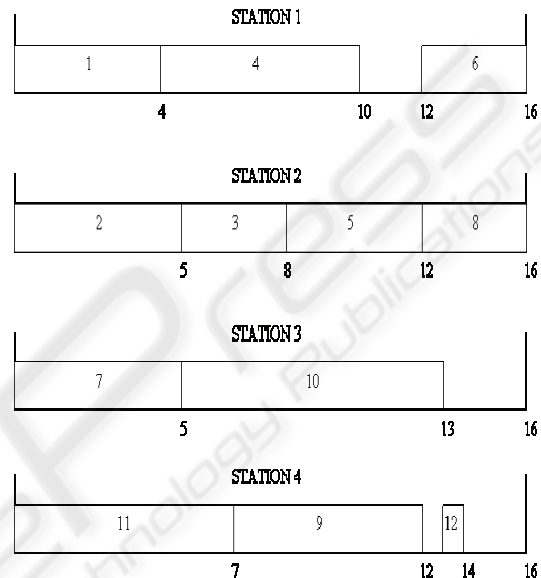


Figure 3: Results for the example problem.

As far as smoothness index and line efficiency are concerned, its estimation, on contrary to LT, is performed without any change to original version. These criterions simply refer to each individual station, despite of parallel character of the method. But for more detailed information about the balance of right or left side of the assembly line additional measures will be proposed:

Smoothness index of the left side

$$SI_L = \sqrt{\sum_{i=1}^K (ST_{maxL} - ST_{iL})^2} \quad (5)$$

where:

SI<sub>L</sub>- smoothness index of the left side of two-sided line  
ST<sub>maxL</sub>- maximum of duration time of left allocated stations  
ST<sub>iL</sub>- duration time of i-th left allocated station

Smoothness index of the right side

$$SI_R = \sqrt{\sum_{i=1}^K (ST_{maxR} - ST_{iR})^2} \quad (6)$$

where:

SI<sub>R</sub>- smoothness index of the right side of two-sided line

ST<sub>maxR</sub>- maximum of duration time of right allocated stations

ST<sub>iR</sub>- duration time of i-th right allocated station

Table 1: Numerical results.

Name	Value
LE	84,38%
LT	30
SI	4,69
SI <sub>R</sub>	2
SI <sub>L</sub>	3

The numerical results of different measures in Table 1 are given. The value of line efficiency is acceptable, smoothness indexes of the right and left side of the line show which part of the assembly line is better balanced. The smoothness index SI informs about balance of the whole line. It is possible to compare the two-sided line balance with single assembly line balance and to consider the influence of position restrictions (L,R or E).

Next it will be consider a small example presented in Figure 4.

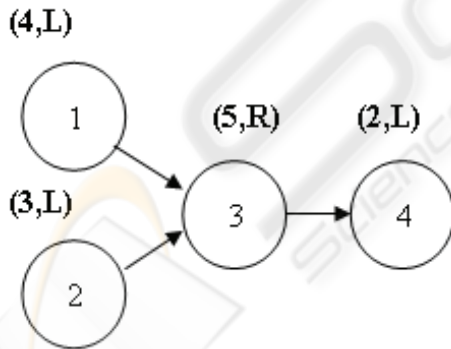


Figure 4: Precedence graph (4 tasks, c=10).

In this point, it's worth to mention about a special case, when mated-station includes instead of two stations, just one. Such a situation takes place, where one station is loaded to a certain point that not allows for assigning any more tasks for this part of the line. As the result, one station stays empty. Balance of this case is presented in Figure 5.

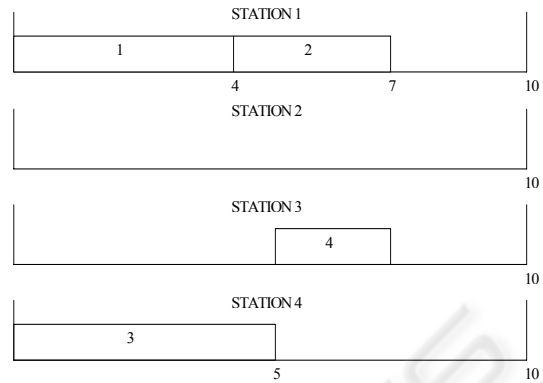


Figure 5: Balance of two-sided line (N=4, c=10).

In this case we got an assembly line which is a structure of incomplete two-sided assembly line. It is possible to estimate the balanced line in two ways: as a single line with parallel stations or incomplete two-sided line.

In the first case we obtain:

$$K = 3$$

$$LE = 46,67\%$$

$$SI = 9,9$$

Considering this case as two-sided line we get:

$$K = 4$$

$$LE = 35\%$$

$$SI_R = 11,18$$

$$SI_L = 8,54$$

$$SI = 14,07$$

As we can see there are some differences in final measurements of the balanced line. The reason is that using heuristic methods we design two-sided assembly line. These kinds of heuristics are very sensitive to cycle time value. Some final balances for different value of cycle time for an example from Figure 2 in Table 2 are shown.

Table 2: Final results of different measures (c = var).

c	K	LT	SI	LE
14	6	37	15,81	66,67%
15	6	39	17,66	62,22%
16	4	30	4,69	84,38%
17	6	43	22,05	54,90%
18	4	32	4,69	77,70%

## 6 CONCLUSIONS

Two-sided assembly lines become more popular in last time. Therefore it is obvious to consider this structure using different methods. In this paper a heuristic approach was discussed. Two-sided assembly line structure is very sensitive to changes of cycle time values. It is possible very often to get incomplete structure of the two-sided assembly line (some stations are missing) in final result. We can use different measures for comparing the solutions (line time, line efficiency, smoothness index). Author proposes additionally two measures: smoothness index of the left side ( $SI_L$ ) and smoothness index of the right side ( $SI_R$ ) of the two-sided assembly line structure. These measurements allow to get more knowledge about allocation of the tasks and about the balance on both sides.

This research was supported by grant of Ministry of Science and Higher Education 3T11Ao2229 in 2005-2008.

## REFERENCES

- Bartholdi J.J., 1993. Balancing two-sided assembly lines: a case study, *International Journal of Production Research*, 23, 403-421
- Baybars, I., 1986. A survey of exact algorithms for simple assembly line balancing problem, *Management Science*, 32, 11-17
- Eral, Erdal, Sarin S.C., 1998. A survey of the assembly line balancing procedures, *Production Planning and Control*, 9, 34-42
- Forteca D.J., Guest C.L., Elam M., Karr C.L., 2005. A fuzzy logic approach to assembly line balancing, *Mathare & Soft Computing*, 57-74
- Gutjahr, A.L., Neumauser G.L., 1964, An algorithm for the balancing problem, *Management Science*, 11, 23-35
- Halgeson W. B., Birnie D. P., 1961. Assembly line balancing using the ranked positional weighting technique, *Journal of Industrial Engineering*, 12, 18-27
- Kao, E.P.C., 1976. A preference order dynamic program for stochastic assembly line balancing, *Management Science*, 22, 19-24
- Lee, T.O., Kim Y., Kim Y.K., 2005. Two-sided assembly line balancing to maximize work relatedness and slackness, *Computers & Industrial Engineering*, 40, 273-292
- Salveson, M.E., 1955. The assembly line balancing problem, *Journal of Industrial Engineering*, 62-69
- Scholl, A., 1998. Balancing and sequencing of assembly line, *Physica- Verlag*
- Sury, R.J., 1971. Aspects of assembly line balancing, *International Journal of Production Research*, 9, 8-14