

SOLVING THE RCPSP WITH AN EVOLUTIONARY ALGORITHM BASED ON INSTANCE INFORMATION

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Abstract: The Resource Constrained Project Scheduling Problem (RCPSP) is NP-hard thus justifying the use metaheuristics for its solution. This paper presents an evolutionary algorithm developed for the RCPSP problem. This evolutionary algorithm uses an alphabet based on random keys that makes easier its implementation while solving combinatorial optimization problems. Random keys allow the use of conventional genetic operators, what makes easier the adaptation of the evolutionary algorithm to new problems. To improve the method's performance, this evolutionary algorithm uses an initial population that is generated considering the information available for the instance. This paper studies the impact of using that information in the initial population. The computational experiments presented compare two types of initial population - the conventional one (generated randomly) and this new approach that considers the information of the instance.

1 INTRODUCTION

The Resource Constrained Project Scheduling Problem (RCPSP) is a classic project scheduling problem, and belongs to the set of Combinatorial Optimization (CO) problems that are very hard to solve and therefore require heuristic procedures. The use of exact methods to solve this type of CO problems is limited to small size instances. The RCPSP problem is a generalization of the production-specific Job-Shop Scheduling Problem (JSSP). According to (Zhang et al., 2007), the Branch and Bound methods for JSSP do not solve instances larger than 250 operations within a reasonable time. As stated in (Liu et al., 2008), in practical manufacturing environments, the scale of job shop scheduling problems could be much larger - in some big textile factories, the number of jobs can be as much as 1,000.

Heuristic methods became very popular and have gained success in solving scheduling problems. For the last twenty years, a huge quantity of papers have been published presenting several metaheuristic methods. From Simulated Annealing to Particle Swarm Optimization (Lian et al., 2006), there are several variants of the same class of method. A very popular method among researchers is the Evolutionary Algorithms (EA).

Vaessens et al. (1996) presented the Genetic Algorithms as the less effective metaheuristic to solve the JSSP. A possibility to increase the efficiency and the effectiveness on an algorithm is to include in the algorithm specific knowledge of the problem. Several works include some specific local search for the JSSP that is based in the critical path of a disjunctive graph. This work presents a strategy to improve the effectiveness and efficiency of an Evolutionary Algorithm to solve the RCPSP. Once verified the difficulty to solve the RCPSP changes from instance to instance, a procedure can be implemented to get knowledge from the instance and transfer it to the EA's initial population. The developed EA is based in a previous one, developed by us, to the JSSP (Oliveira et al., 2010) according to the similarities between both problems. Above all of this, the strategy to enhance the effectiveness was tested in the JSSP by Oliveira et al. (2010) and in a very particular case it improves the makespan in about 8%.

The paper is organized as follows: an introductory section defines the RCPSP problem and its representation with the use of an AoN graph; a central section describes the methodology and a numerical example that supports the explanation of the constructive algorithm; finally, some conclusions and a discussion on future work are presented.

2 RCPSP

A project can be represented as a network of n activities in a graph, where exist links between pairs of activities, representing processing precedence between them. The order to process the set of activities must respect the precedence set. There are two ways to represent a project: the activity-on-arc (AoA) and the activity-on-node (AoN).

In the AoA mode, the set of nodes represents the "events" (start/end processing) and the set of arcs represents the "activities." In the AoN scheme, the set of nodes represents the "activities" and the set of arcs represents the precedence between the "activities." In general, each activity requires the simultaneous use of several resources (Demeulemeester et al., 2003). In this paper, the AoN scheme has been adopted since it allows a direct correspondence with the disjunctive graph that is usually used to represent the JSSP. Figure 1 shows a project network with 10 activities, represented in a AoN scheme (Ranjbar and Kianfar, 2009).

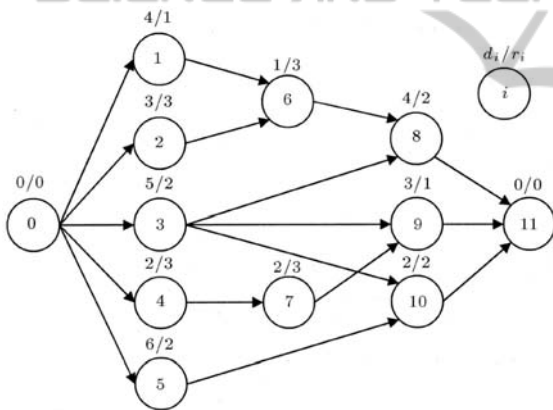


Figure 1: AoN project network.

For a given r renewable resources R_1, \dots, R_r , a constant amount of b_k units of resource R_k is available at any time. Activity j must be processed for d_j time units, where preemption is not allowed. During this period of time, a constant amount of r_{jk} units of resource R_k is occupied. The objective is to determine the starting times of S_j for the activities $j = 1, \dots, n$ in a way that:

- i) at each time t , the total resource demanded is less than or equal to the resource availability for each resource type;
- ii) the given precedence constraints are fulfilled;
- iii) the makespan $C_{\max} = \max_{j=1}^n C_j$, where $C_j = S_j + d_j$, is minimized.

Figure 1 shows also 2 dummy activities, node 0

and node 11 which represent the starting time of the project (node 0) and the project conclusion time (node 11). Their duration is null. Each node shows the duration of the activity and the amount of renewable resources. In this project instance it is considered only one renewable resource. The set of arcs represent the precedence set of the project. Activity 6 can only be started after activities 1 and 2 are completed.

The RCPSP is a highly complex optimization problem due to its combinatorial nature. This problem belongs to the NP-hard class (Brucker et al., 1998), and so an efficient algorithm for obtaining exact solutions is unknown.

The research on the RCPSP has been extensive along the years and it involves exact algorithms and heuristic methodology. Generally, the exact methods are based on branch and bound algorithm, where the branch strategies are related to the choice of activities in the disjunctive graph that represents the project. Also, in this field of research, the continuous development of new and better lower bounds for the optimal solution is pointed out, what is a fundamental auxiliary information for the heuristic methods.

The state-of-the-art RCPSP performed by Hartmann and Kolisch (2000) and updated later (Kolisch and Hartmann, 2006) shows a bigger number of publications on heuristics than on exact methods. In this study, the authors explain that, in the specialized bibliography, there are several heuristic methodologies to deal with RCPSP. These methods can be classified as "single pass" (over activities set) or "multiple pass", based in priority dispatch rules to construct the solution, local iterative improvement and meta-heuristics.

In the metaheuristics set, several approaches have arisen, such as: evolution strategies, taboo search, ant colony system, variable neighbourhood search and simulated annealing. Some of those metaheuristics include hybridization with local search procedures and with single pass or multiple pass. An example of hybridization is given by Ranjbar and Kianfar (2009). They presented a hybrid metaheuristic algorithm for solving the RCPSP, which contains a scatter search skeleton and uses a special solution combination method for making children.

The process to construct a solution is a fundamental component in the heuristic methodology. With RCPSP in particular, this component is called Schedule Generation Scheme (SGS). For the RCPSP, two SGS models exist: the serial model and the parallel model. It is possible to find in the literature heuristic procedures that use

either one scheme or the other, but it is also possible to find some methodologies that use both, or even a combination of the two models.

Analyzing the computational results presented in the literature performed with test instances with 30, 60 and 120 activities, it is possible to point out that the better approximation methodology is the population-based heuristic with representation of list of activities and serial SGS. The parallel SGS seems to be better when dealing with large instances.

It seems that further research for approach methods will concentrate on the development of algorithms which integrate the Forward-Backward Improvement strategy (FBI), since four of the best six heuristics already use this local search strategy. The FBI strategy has two phases: the first phase gives a feasible schedule using a priority dispatch rule to sequence the activities in a bidirectional way; the second phase tries to improve the resources assignment performing successive bidirectional passes without increasing the project duration obtained in the first phase.

The performance of a heuristic can be improved by integrating specific knowledge of the problem. With the RCPSP, the integration of a constructive algorithm in an Evolutionary Algorithm (EA) constructs only active schedules. This strategy is implemented similarly to that performed for the particular case of the Job Shop Scheduling Problem.

3 EVOLUTIONARY ALGORITHM

This work adopts a method based on Evolutionary Algorithms. EA can be viewed as a metaheuristic search techniques that imitates the principle of evolution in nature i.e., EAs use a set of competing potential solutions of the problem which evolve according to rules of selection and transformation.

It is possible to briefly describe the functioning of an EA in the following way. An EA proceeds in iterations called generations. In each generation, a set of solutions called population is generated. Usually, the first population is generated randomly. A new population is generated from a current population using selection and transformation operators. The members of a population are called individuals. The fitness-function assigns a value called fitness to each individual. Given an individual and its fitness, a selection operator decides if an individual from the current population is used as an input of a transformation operator. A transformation operator creates a new element of a new population from an arbitrary number of elements of the current

population. It is expected that during the search process, increasingly better solutions are found. To achieve this goal, it is necessary that the transformation operators focus attention on high fitness. Very common transformation operators are reproduction (copying an individual to the new population unaltered), mutation (an individual is changed in a random fashion) and crossover (where parts of two individuals are combined).

The simplicity of an EA to model more complex problems and its easy integration with other optimization methods were factors that were considered for its choice. One advantage of this EA is the portability to other problems. The algorithm proposed was conceived to solve the RCPSP, but the method can be adapted to solve other variants or extensions of the RCPSP, such as the RCMPSP. In fact, the EA was already adapted to solve other type of problems in the Project Planning field (Silva et al., 2010). When implementing an appropriate algorithm to construct solutions based on chromosome, the EA can also solve other scheduling algorithms like job shop (Oliveira et al., 2010). Furthermore, random keys chromosome (Bean, 1994) also enables to apply the algorithm in other type of problems once conventional genetic operators, which are problem independent, could be used.

As Gonçalves et al. (2005) state, the important feature of random keys is that all offspring formed by crossover are feasible solutions, when it is used as a constructive procedure based on the available operations to schedule and the priority is given by the random key allele. Through the dynamics of the genetic algorithm, the system learns the relationship between random key vectors and solutions with good objective function values.

A chromosome represents a solution to the problem and is encoded as a vector of random keys (random numbers). In this work, according to Cheng et al. (1996), the problem representation is indeed a mix from priority rule-based representation and random keys representation. An excellent classification of heuristics algorithms for the RCPSP was provided by Kolisch and Hartmann (1999). They discussed several methods and alternatives to generate different type of solutions for the RCPSP and, in particular, they provided detailed algorithms to generate active schedules. In the meantime, in this EA the solutions are decoded / constructed by an algorithm, that is based on Giffler and Thompson's algorithm (Giffler and Thompson, 1960) for the JSSP. While the Giffler and Thompson's algorithm can generate all the active

plans for the JSSP, the adapted constructive algorithm only generates a plan according to the chromosome. Similar to the JSSP, the solutions of RCPSP can be classified on three sets: Semi-active, Active, and Non-delay schedules (Sprecher et al., 1995).

As advantages of this strategy, we have pointed the minor dimension of solution space, which includes the optimum solution and the fact that it does not produce impossible or disinteresting solutions from the optimization point of view. On the other hand, since the dimensions between the representation space and the solution space are very different, this option can represent a problem because different chromosomes can represent the same solution.

Using the adapted Giffler and Thompson's algorithm to construct a solution based on the priority values of each activity, a Serial Schedule Generation Scheme is adopted. It has been shown by Kolisch (1996) that a serial scheme yields better results when a large number of schedules is computed for one project instance. Attending Hartmann (1998), the parallel scheme might exclude all optimal solutions from the search space, while the search space of the serial scheme always contains an optimal schedule.

The constructive algorithm has n stages and in each stage an activity is scheduled. To assist the algorithm's presentation, consider the following notation existing in stage t :

- P_t - the partial schedule of the $(t-1)$ scheduled activities;
- S_t - the set of activities schedulable at stage t , i.e. all the activities that must precede those in S_t are in P_t ;
- Q_t - the set of activities queued at stage t , i.e. all the activities that are not in S_t nor in P_t ;
- σ_k - the earliest time that activity a_k in S_t could be started;
- ϕ_k - the earliest time that activity a_k in S_t could be finished, that is $\phi_k = \sigma_k + p_k$;
- M^* - the set of resources used by a_k in S_t which has $\phi^* = \min_{a_k \in S_t} \{\phi_k\}$;
- S_t^* - the conflict set formed by a_k in S_t which use at least one resource of M^* and $\sigma_j < \phi^*$;
- a_j^* - the selected activity to be scheduled at stage t .

Table 1 presents the constructive algorithm that is used to build a solution.

In *Step 3*, instead of using a priority dispatching rule, the information given by the chromosome is used. If the maximum allele value is equal for two or more operations, one is chosen randomly.

Table 1: A constructive algorithm.

<i>Step 1</i>	Let $t=1$ with P_1 being null. S_1 will be the set of all activities with no predecessors.
<i>Step 2</i>	Find $\phi^* = \min_{a_k \in S_t} \{\phi_k\}$ and identify M^* . Form S_t^* .
<i>Step 3</i>	Select activity a_j^* in S_t^* , with the greatest allele value.
<i>Step 4</i>	Move to next stage by <ol style="list-style-type: none"> (1) adding a_j^* to P_t, so creating P_{t+1}; (2) removing a_j^* from the list of predecessors of activities in Q_t deleting a_j^* from S_t and creating S_{t+1} by adding to S_t the activities that a_j^* as the only predecessor; (3) deleting a_j^* from S_t and creating S_{t+1} by adding to S_t the activities with no predecessors; (4) incrementing t by 1.
<i>Step 5</i>	If there are any activities left unscheduled ($t < N$), go to <i>Step 2</i> . Otherwise, stop.

3.1 First Generation

In the RCPSP it is possible to calculate for each activity the remaining time for completion of the project. In the project network this time is referred to as "tail". This value represents the longest path from the activity to node $n+1$. It is easy to admit that the activities with large tail must be sequenced first because they could define the makespan. Indeed, there exists the dispatching rule that is used in practice JSSP which sequences operations in first place that belongs to the job with Most Work Remaining (MWR).

Consider the following example presented in Figure 1, with ten activities. In this instance there is only one constrained renewable resource with 4 available units. Table 2 presents this instance. The activities are numbered sequentially and represented by index j . The duration of each activity is given in the row "duration".

Table 2: The instance data.

Activity	1	2	3	4	5	6	7	8	9	10
duration	4	3	5	2	6	1	2	4	3	2
resources	1	3	2	3	2	3	3	2	1	2
tail	9	8	9	7	8	5	5	4	3	2

The *tail* row shows the remaining time to complete the project including the processing of the activity. For instance, the tail of activity 3 is 9 and is equal to the duration of activities 3, 8 and, 11, and that is respectively 5, 4, and 0 (Figure 1). It correspond to the longest path from node 3 to node 11.

Applying the MWR rule to this instance the schedule is given in Figure 2. For this small example, the optimal solution is presented in Figure 3.

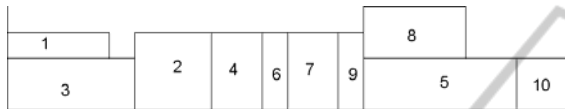


Figure 2: Gantt chart of MWR solution.



Figure 3: Gantt chart of optimal solution.

In the chosen random key representation, a chromosome is a vector of N genes that are real numbers between 0 and 1. The representation has one gene for each activity. Table 3 presents two chromosomes for a RCPSW of ten activities. The second chromosome (*chrms2*) was generated randomly, while the first chromosome (*chrms1*) represents the information that is obtained dividing in the example above the tail value per maximum tail value plus one (9+1 in the example). Applying the constructive algorithm and considering the chromosome *chrms1*, the solution of Figure 2 is obtained. Attending to this property we implement a procedure to generate the first population that includes some knowledge about the instance in the form of the tail of each activity.

Table 3: Chromosomes for a RCPSW10.

Activity	1	2	3	4	5	6	7	8	9	10
chrms1	0.90	0.80	0.90	0.70	0.80	0.50	0.50	0.40	0.30	0.20
chrms2	0.30	0.30	0.70	0.80	0.30	0.40	0.20	0.70	0.10	0.20

To generate *L* individuals in a population the following expression at (1) is used to generate the allele value for each gene *j*:

$$allele_j = \frac{randbetween(tail_j; tail_j + i \times GAP)}{\max_j(tail) + i \times GAP} \quad (1)$$

3.2 Considerations About GAP

The *GAP* value allows obtaining at the end of population a quite random chromosome. The higher the *GAP*, the more random is the population. Lower *GAP* gives a population close to the MWR dispatch rule.

For a population of 50 individuals, the graphic presented in Figure 4 presents some statistics of the allele values for each gene (10 activities / genes). The average value, maximum value, minimum value and also the tail value of each operation is presented.

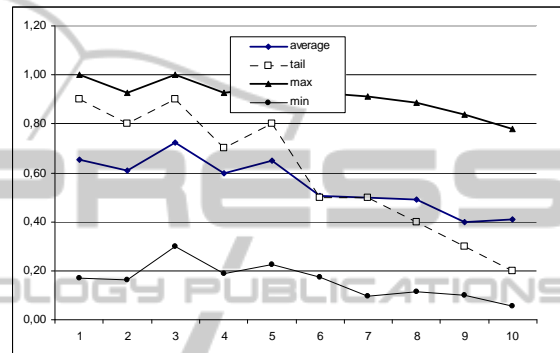


Figure 4: Statistics of allele's values.

Figures 5 and 6 show the *GAP* effect in the statistics of the allele values for the 10 activities. Figure 5 shows the statistics when a small *GAP* is used. The values are very close to the tail of each operation.

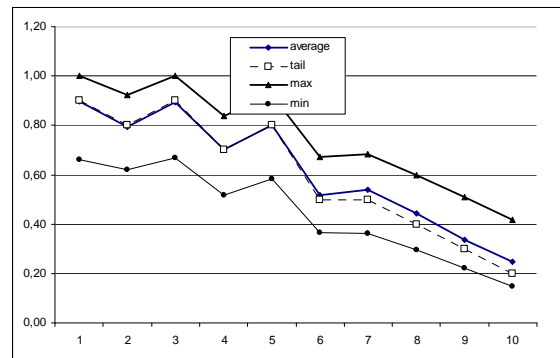


Figure 5: Statistics of allele's values with small GAP.

Figure 6 shows the statistics for a high *GAP*. The range of variation for the alleles is very similar for all genes, so the population is "more random".

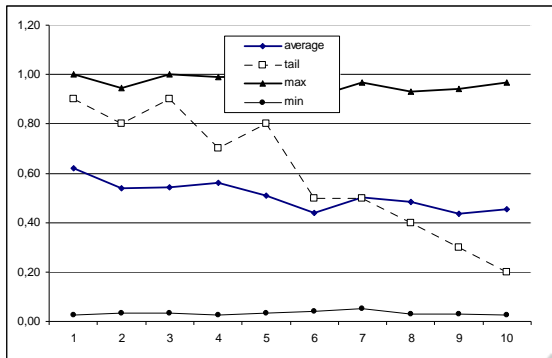


Figure 6: Statistics of allele's values with high GAP.

Figure 7 presents the alleles' values of four genes when using a GAP=2. These genes belong to the activities having different tails values which are respectively 9, 7, 5, and 2. It is possible to verify that for the first chromosomes the alleles' values translate the priority give by the tail of each operation. It is also possible to verify that the last individuals have quite random alleles. It is possible to see that in some individuals the activity 10 (Gene10) and activity 6 (Gene6) have a priority greater than the activity 1 (Gene1). This situation allows a first sequence of these activities before activity 1, which is the activity with the greatest tail.

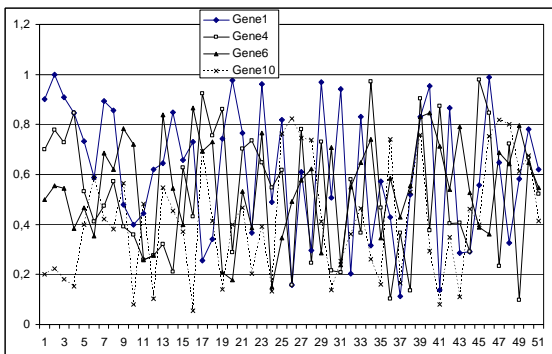


Figure 7: The allele generation.

Figures 8 and 9 show the effect of GAP on the allele generation for the same genes. Figure 8 shows the alleles values when used as small GAP. The values are very close to the tail of each operation.

Figure 9 shows the statistics for a high GAP. The range of variation for the alleles is very similar for all genes, so these allow "any" sequence to schedule these four activities.

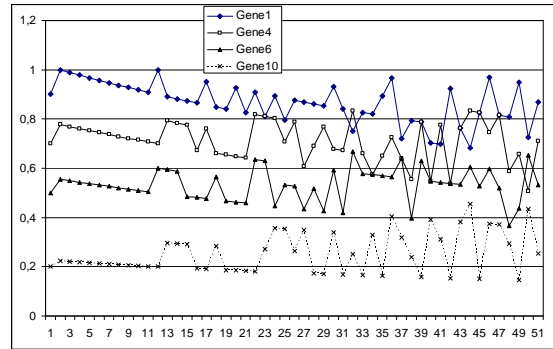


Figure 8: The allele generation with small GAP.

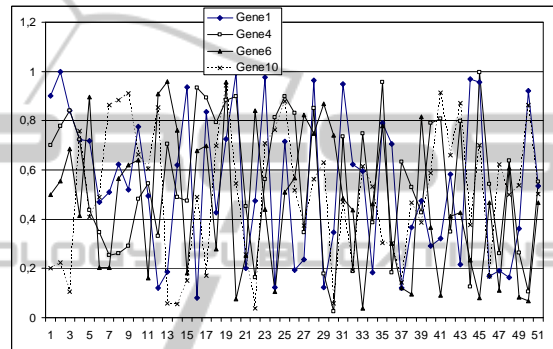


Figure 9: The allele generation with high GAP.

3.3 The Algorithm Structure

The genetic algorithm has a very simple structure and can be represented in Table 4. It begins with population generation and her evaluation. Attending to the fitness of the chromosomes the individuals are selected to be parents. The crossover is applied and it generates a new temporary population that also is evaluated. Comparing the fitness of the new elements and of their progenitors the former population is updated.

Table 4: Genetic algorithm.

```

begin
P ← GenerateInitialPopulation()
Evaluate(P)
while termination conditions not meet do
    P' ← Recombine(P) //UX
    Evaluate(P')
    P ← Select(P ∪ P')
end while
    
```

The Uniform Crossover (UX) is used this work. This genetic operator uses a new sequence of random numbers and swaps both progenitors' alleles if the random key is greater than a prefixed value.

Table 5 illustrates the UX's application on two parents (prnt1, prnt2), and swaps alleles if the random key is greater or equal than 0.75.

Table 5: The UX crossover.

<i>i</i>	1	2	3	4	5	6	7	8	9	10
prnt1	0.89	0.48	0.24	0.03	0.41	0.11	0.24	0.12	0.33	0.30
prnt2	0.83	0.41	0.40	0.04	0.29	0.35	0.38	0.01	0.42	0.32
randkey	0.64	0.72	0.75	0.83	0.26	0.56	0.28	0.31	0.09	0.11
dscndt1	0.89	0.48	0.40	0.04	0.41	0.11	0.24	0.12	0.33	0.30
dscndt2	0.83	0.41	0.24	0.03	0.29	0.35	0.38	0.01	0.42	0.32

The genes 3 and 4 are changed and it originates two descendants (dscndt1, dscndt2). Descendant 1 is similar to parent 1, because it has about 75% of genes of this parent.

4 EXPERIMENTS

The Genetic Algorithm was coded in C++ and the computer used was a PC Pentium IV 3GHz processor with 1024 Mbytes of RAM., with Linux operative system.

The problem data set was taken from PSPLIB datasets (Kolisch and Sprecher, 1997). The dataset consists of four test sets J30, J60, J90 and J120 that contain problem instances of 30, 60, 90 and 120 activities, respectively, and have been constructed by the instance generator ProGen. In this preliminary study we present results for 4 instances with 30 activities: J301_1, J301_2, J301_3 and J301_4. The optimal value for these instances are respectively 43, 47, 47 and 62.

Table 6 resumes the computational experiments. For each instance were made five runs considering the tail information on the initial population, and also were made five runs with an initial population generated randomly. In each run the algorithm performs 5000 iterations.

Table 6 shows the average value and the best value of five runs. Each column shows the values in different iterations of the algorithm. In the computation experiments it was only used a value for the GAP.

Globally in all experiments the use of tails' information of the instance produces in average 1.8% better results. This advantage is more evident in the initial iterations, so for some instances it was obtained solutions about 5.5% better. In terms of best solution the use of tails' information is relevant on the initial iterations. In some instances the use of tails information produces a solution about 3.6% better than using a random initial population.

Table 6: Computational experiments.

Iteration	J301_1_Random		J301_1_tails	
	Average	Best	Average	Best
1	76.2	73	76.2	75
50	61.8	60	58.4	58
100	60.8	59	58.4	58
500	57.0	55	54.0	54
1000	56.0	54	54.0	54
2000	55.0	53	53.0	53
5000	53.4	53	53.0	53
Iteration	J301_2_Random		J301_2_tails	
	Average	Best	Average	Best
1	89.2	84	90.4	90
50	84.2	84	82.6	81
100	84.0	84	82.6	81
500	80.8	80	78.4	78
1000	80.4	80	78.4	78
2000	78.8	77	78.4	78
5000	78.4	77	78.4	78
Iteration	J301_3_Random		J301_3_tails	
	Average	Best	Average	Best
1	93.0	91	89.6	89
50	81.6	81	80.6	79
100	81.0	81	79.8	79
500	78.6	77	79.0	78
1000	76.6	75	76.0	76
2000	76.4	75	76.0	76
5000	75.4	75	76.0	76
Iteration	J301_4_Random		J301_4_tails	
	Average	Best	Average	Best
1	152.6	146	149.6	147
50	148.0	146	144.4	143
100	145.4	142	144.4	143
500	140.2	138	136.0	136
1000	139.6	138	136.0	136
2000	139.0	138	136.0	136
5000	137.4	136	136.0	136

5 CONCLUSIONS

This paper presents a new method for generating the initial population for an evolutionary algorithm that solves the Resource Constrained Project Scheduling Problem (RCPSP). This proposed method, through known information for each instance, establishes different priorities to the activities based on the remaining processing time to finish the project, thus generating an initial population. The algorithm uses a representation based on random keys.

Results are then compared with the conventional random generation and some promising conclusions arise. In fact the experimental computation shows better results when the information of the instance is used, especially on initial iterations.

Definitely this could be a very simple way to increase effectiveness of the algorithm. Furthermore the use of instance information could also be parameterized, obtaining levels of convergence of the algorithm depending on available CPU time.

Since the random keys representation has no Lemarkin property, it is our intention to apply the

same strategy on initial population generation in other types of representation for the RCPSP, namely the permutation code.

Further work would also consist on applying this method to other problems. The Resource Constrained Multi-Project Scheduling Problem (RCMPSP), as a generalization of the RCPSP, would be the next candidate for evaluating the advantages of the approach presented above.

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