

Force Directed Flow Map Layout

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Abstract: A flow map is a thematic map that is been used to emphasize the spatial pattern of one or more geographic attributes. Although this kind of thematic maps is often drawn by hand, a few automatic computer algorithms exist. In this research paper, we proposed a novel algorithm for the automatic generation of flow maps that is theoretically grounded on physics' laws to describe the motion and force of attraction or repulsion between points. Properties associated to these laws are then used to merge different flows, as well as for the improvement of the maps' visual quality. Finally, we evaluate our work by generating a set of flow maps and by doing a comparison with flow maps produced by existing algorithms.

1 INTRODUCTION

Geographic flows can represent the movement of tangible objects (e.g. people, bank notes, and goods) in geographical space, but also of intangible objects (e.g. energy, ideas, and reputation). The functional definition of 'flow' is 'the continuous movement of objects in one direction' (Cambridge, 2013). The quantification of movement (e.g. the number of moving objects or the number of transitions) within a flow is called 'flow magnitude' (Andrienko, et al., 2007). The most frequent questions associated with this kind of data are the following: where does the flow start? Where does the flow arrive? Which is the magnitude of the flow? And which is the magnitude of flows that share the same destination?

An answer to these questions can be easily deduced using visualization techniques like flow maps. A flow map is a thematic map, which gives

emphasis to the spatial pattern of one or more geographic attributes (Slocum, et al., 2009). In particular, it shows the spatial distribution of univariate geographic phenomena. Additional properties of flow maps are: the flow magnitude is represented as the width of each flow line, and the sum of all flow line branches should add up to the width of the flow line (Dent, 1990).

The first flow map was created by a cartographer named Henry Drury Harness, in 1837, by showing the transportation of passengers in Ireland. Since then, this kind of thematic map is been used to represent flow data. The success of this visualization technique is due to its simple design and intuitive understanding. Several years ago, a French cartographer called Charles Joseph Minard introduced one of the most popular flow maps techniques which, for example, he used to represent the movement of travelers on the principal railroads of Europe in 1865, as well as the international

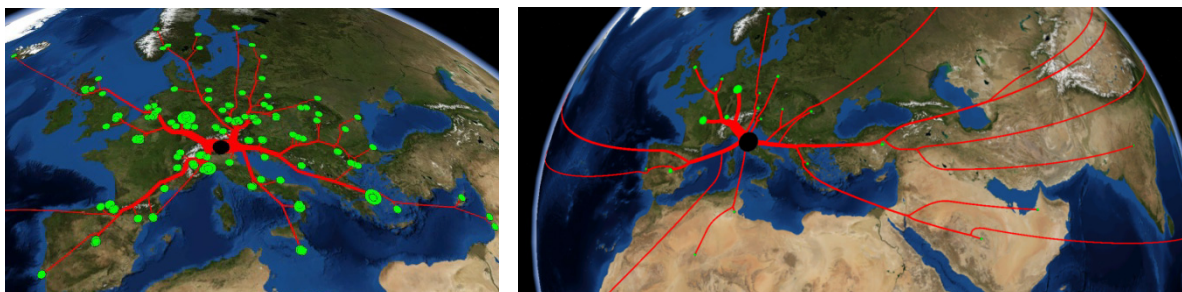


Figure 1: Flow maps: (left) Connections between Fondazione Graphitech and its partners (research centers, university, companies). (right) Main exports of Italian goods in all over the world.

distribution of French wines, cotton and coal (Friendly, 2000). Many of his works are still available online (datavis.ca, 2013).

In flow maps, aggregation techniques can be of valuable use because they reduce the visual clutter (e.g. each flow is easier to trace). Additionally, they enable the visualization of the magnitude of the flows that share common destinations.

2 RELATED WORK

In this section we summarize the work done to automatically generate flow maps, and then we provide an overview of the force directed algorithms in the field of graph drawing domain.

2.1 Algorithms for Automatic Generation of Flow Maps

In this subsection we describe existent algorithms for the automatic generation of flow maps.

Tobler in 1987 (Tobler, 1987) developed a computer program called FlowMapper (Tools, 2001) used to visualize migration maps; this mapper produces generic maps without any optimization with respect to the visualization of the data. For each flow, it generates an arrow on the map with a width varying accordingly to its magnitude. As the number of links increases, it becomes increasingly difficult to represent new flows without creating a visual clutter (see Figure 2a). Occlusion of trajectories produces maps that are difficult to interpret unless some form of generalization is applied.

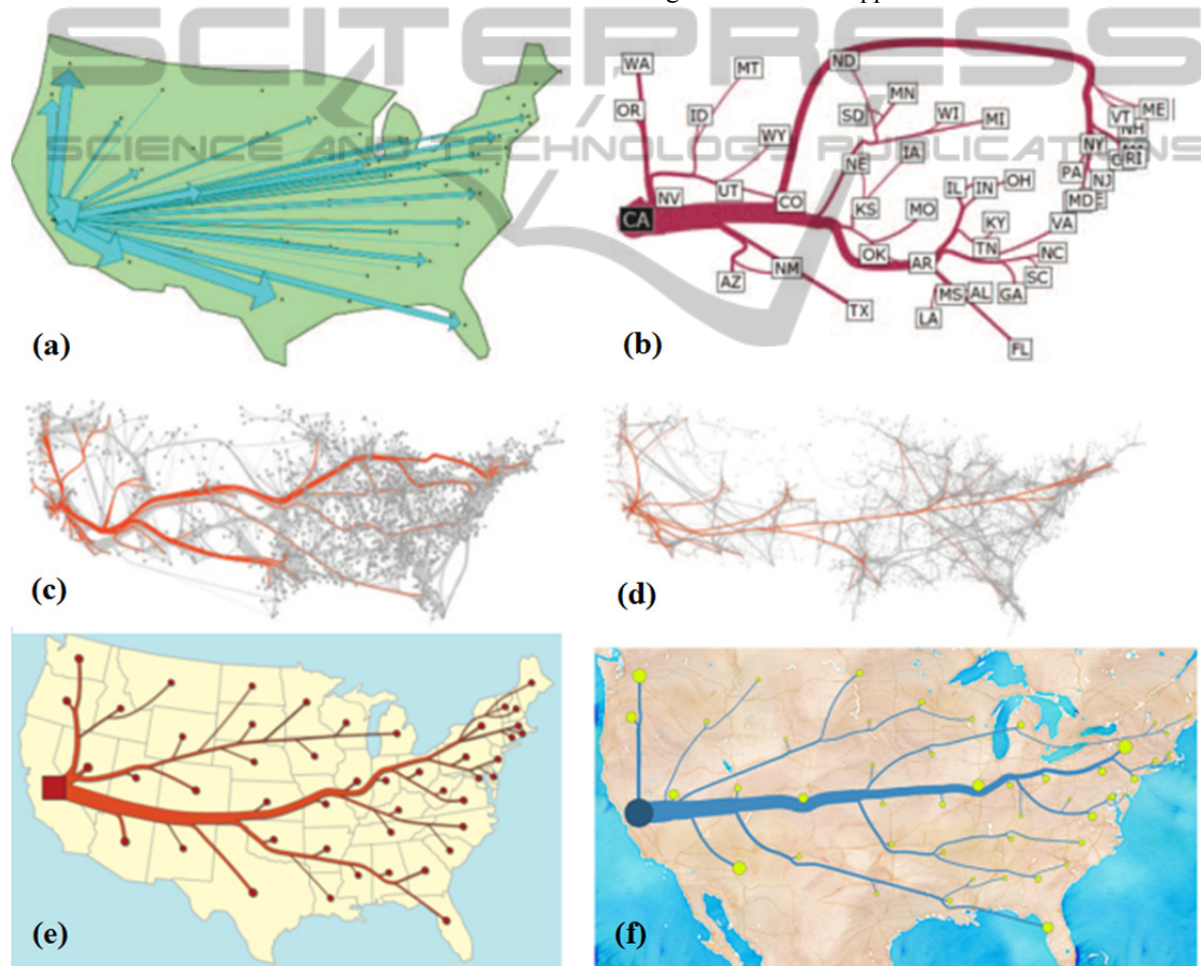


Figure 2: Different maps representing the migration from California in 1995-2000. (a) the first software that automatically generate flow maps (Tools, 2001), (b) algorithm taking into account aggregation of flows (Phan, et al., 2005), (c) (d) the force directed algorithms in graph drawing domain (Cui, et al., 2008), (Holten and Van Wijk, 2009), (e) the algorithm based on spiral tree (Verbeek, et al., 2011), (f) the force directed algorithm described in this paper.

The main aspect of a flow map is the aggregation of the flows. In 2005 Phan et al. (Phan, et al., 2005) introduced a method to reduce the visual clutter by merging the flows (see Figure 2b). The author used a binary hierarchical clustering formulating the layout in a simple and recursive manner. The advantage of this algorithm is its $O(n^2)$ time complexity, where n is the total number of nodes used in the algorithm. However this algorithm has a few limitations: during the first step the nodes are moved if their proximity is too small. Hence, it might lose the geographical reference associated to each node. Furthermore, if there are too many destination nodes in a small area, by forcing binary splits introduces too many extra routing nodes, which then leads to clutter.

Verbeek et al. (Verbeek, et al., 2011) describes a method to overcome the aforementioned limitations through the use of spiral trees. The authors used spiral spline (Buchin, et al., 2011), which have $O(n \log n)$ complexity, to generate maps that are crossing-free. Another key property concerns the weights of the leaves. If there are two flow maps with the same origin and destinations but with different magnitudes, then the layout of the flow map will be represented differently. Moreover the flow tree produced is constrained to avoid crossing its own nodes, as well as user-specified obstacles (see Figure 2e). In order to have a high-quality map, a quality function that takes into account the obstacles, the smoothness, the angles and the straightness, has to be minimized. The time required to perform a flow map is mentioned in the paper as a “couple of minutes”. Two limitations of this algorithm are the complexity in the construction of the tree structure and the non-intuitive body cost function used to improve the aesthetic results.

2.2 Force-directed Techniques

The reduction of visual clutter has been studied extensively in the graph drawing domain. A well-studied class of approaches to overcome this challenge is composed by force-directed algorithms. In few words, in a force-directed algorithm the graph is represented as a physical system of particles with forces acting between them. At each iteration, the energy of the system - defined as the sum of all the forces that affects the particles - changes. The algorithm halts when the local minimum of the energy is found.

Eades et al. (Eades, 1984) introduced the idea of replacing vertices by hinges and edges by springs. Vertices are placed in a given initial positions, and springs move the vertices to a minimal energy state

(i.e. the springs are compressed or extended as little as possible). A few years later, similar methods were introduced as an extension to this idea.

Kamada and Kawai's algorithm (Kamada and Kawai, 1989) added the concept of an ideal distance: the ideal distance between two vertices is proportional to the length of the shortest path between them.

Fruchterman et al. (Fruchterman and Reingold, 1991) presented a technique to draw general undirected graphs according to some aesthetic criteria, such as the distribution of the vertices in the frame and the minimization of edge crossings. In this case the forces are applied to each node of the graph. Dwyer et al. proposed to add edge routing to force-directed layouts (Dwyer, et al., 2007).

Cui et al. (Cui, et al., 2008) used a control mesh that reflects the underlying control pattern reducing the visual clutter (see Figure 2c).

In the latter case Holten et al. (Holten and Van Wijk, 2009) presented a force-directed algorithm in which the edges are modeled as flexible springs that can attract each other while node positions remain fixed (see Figure 2d). This algorithm works well in aggregating edges of a generic graph. However the generated graphs do not merge as quickly and smoothly as hand-drawn flow maps.

Our work improves the algorithm proposed by Holten et al. (Holten and Van Wijk, 2009) to aggregate single origin flow tree.

3 SYSTEM DESIGN

The core of our algorithm focalizes on force-directed techniques, following simple and intuitive concepts that can be implemented in a few lines of code and easily extended for possible improvements.

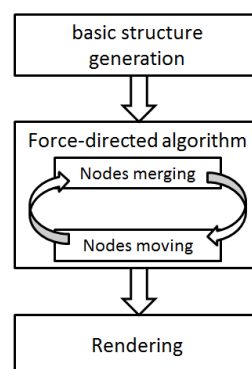


Figure 3: System diagram.

As mentioned in the section 0, algorithms to automatically generate flow maps first create a flow map tree and then use such structure to generate the flow map. In this work, flows are generated directly using a force-directed approach without creating a tree structure a priori. The only structure necessary is composed by the edges that start from the origin and arrive to the target nodes.

This work has the objective of achieving automatic generation of flow-maps, which are created in a natural way, with high visual quality.

In particular we want to satisfy the following aesthetic criteria, which are a common goal of previous works as well (Phan, et al., 2005), (Verbeek, et al., 2011):

1. The possibility of aggregate flows, reducing the visual clutter.
2. The use of smooth curves for aesthetic purposes.
3. The flow magnitude affects the layout of the generated map; straight lines are correlated to nodes with high magnitude.
4. The target nodes are not overlapped with flows.
5. The flow is crossing-free.

In contrast to previous works, this paper presents a novel force-directed algorithm to automatically generate flows maps that satisfy the aforementioned criteria and that is intuitive: imagine a pavement with a set of elastic cords, tied on one side to a common stake, and on the other side to different stakes placed over the pavement. Every cord has a different elasticity. In this example, we attempt to aggregate these cords. Hence, we start from the stake where all cords are tied and we try to bind together near cords; those with higher elasticity are dragged by the ones with less elasticity. However, we have two criteria that are not satisfied: elastic cords do not have smooth appearance and they might collide with stakes. These are two issues of our algorithm that shall be addressed in the following sections.

The implemented method is divided in different steps (see Figure 3).

3.1 Basic Structure Generation

In order to aggregate flows that initially are straight lines (see Figure 6a), we use a subdivision technique to divide each flow line into a set of smaller segments, hence creating new intermediate nodes. However, not all initial flow lines end up with same

number of intermediate nodes. The number of intermediate nodes generated for the flow is proportional to the distance to the target. Each intermediate node has an index, called “node index” that defines the order. Nodes having the same “node index” maintain the same distance from the origin (see Figure 6b). The idea is to aggregate the node with the same “node index”.

At this step each intermediate node has one father node and one child node. The flow magnitude will be assigned to each intermediate node, thus each intermediate node has the flow magnitude of its flow.

Each intermediate node also knows which are the two intermediate nodes with the same node index that belong respectively to the previous and following flow taking into account the clockwise order. We define these two intermediate nodes as “near nodes”.

3.2 Force-directed Algorithm

This phase is divided into two crucial steps, executed every iteration.

3.2.1 Nodes Merging

This phase is responsible for the update of the flow tree that becomes obsolete after merging the nodes, having the same node index, the same father and within a certain distance.

Another mandatory condition is that the two nodes are “near nodes” with each other. In this way crossing between flows cannot occur. This satisfies the criterion 5, for example looking at the Figure 6 the nodes that belong to the flow N can only merge with the nodes of flow M, as well as the nodes of the flow Q. Meanwhile the nodes of the flow M can merge both with the nodes of flow N and Q.

When the merge event occurs, the two nodes will be removed and a new node is created to represent their merge (see Figure 6c). The condition of having the same father is required to maintain the consistence of the tree structure. The new node will inherit the behavior of the two nodes merged. Additionally, its flow magnitude will be the sum of the two node’s magnitudes. Hence, at each iteration two or more intermediate nodes are aggregated, reducing the number of nodes and updating the tree structure. This step satisfies the aesthetic criterion 1. In fact merging the nodes that compose the structure implies the merging of the flows.

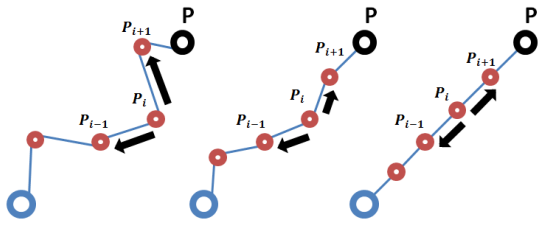


Figure 4: Stress force applied to the intermediate nodes of a flow.

3.2.2 Nodes Moving

In this phase every intermediate node will interact with zero, one or more other intermediate nodes. The interaction is possible using different forces.

Electrostatic Force.

Attractive forces are calculated only between nodes sharing the same index and father. This kind of force is called electrostatic force and it is described by the following equation:

$$F_e(p_i, q_i) = \frac{1}{\|p_i - q_i\|} (\widehat{p_i - q_i}) \quad (1)$$

where p_i is the vector representing the selected node and q_i is the vector representing the interacting node of index i . The \widehat{A} notation means the unit vector of A meanwhile the $\|A\|$ notation means the norm of A . The closer are two nodes and the higher will be the force between them. However this formula takes only in account the distance of the interacting nodes and not their magnitude. Thus, we only apply the force if the condition $m_{q_i} \geq m_{p_i}$ is satisfied. m_{q_i} and m_{p_i} are respectively the magnitude of the node p_i and q_i .

Hence, nodes of smaller magnitude will be attracted by nodes of higher magnitude. Additionally, flows having higher magnitude will be straighter than those one having lower magnitude. Consequently, the aesthetic criterion 3, which states "the flow magnitude affects the layout of the generated map; straight lines are correlated to nodes with high magnitude", is reached.

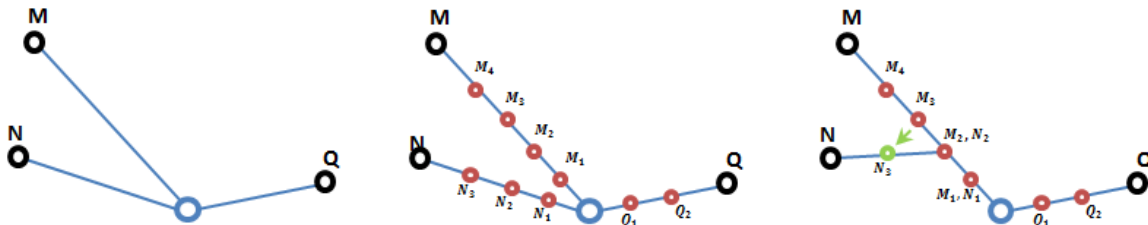


Figure 6: Steps of the algorithm, (a) initial situation, (b) creation of the basic tree structure, (c) aggregation of nodes.

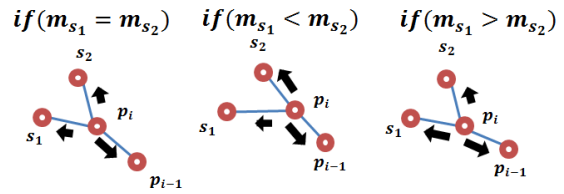


Figure 5: Different cases of stress force taking into account the flow magnitude of child nodes.

Stress Force.

In order to keep a homogeneous position from the father node and from the child node, each node applies a force to these nodes.

$$F_s(p_i) = (p_{i-1} - p_i) + (p_i - p_{i+1}) \quad (2)$$

p_{i-1} and p_{i+1} are respectively the vector representing the father and the vector representing the child of the node p_i . This is the force where each node has one father and one child node. This force tends to offset each node to a straight line formation (see Figure 4). When nodes have two or more children, we apply the following formula:

$$F_s(p_i) = (p_{i-1} - p_i) + \sum_{s \in S_{p_i}} \frac{m_s}{m_{p_i}} (p_i - s) \quad (3)$$

S_{p_i} is the set of the children of the node p_i . The node is moved towards the flow of higher magnitude (see Figure 5). m_{p_i} is the magnitude of the current node, that is, the sum of the magnitude of all children. m_s is the magnitude of one child node. In order to obtain smooth lines the stress force is applied only if the force is greater than a threshold t : $F_s(p_i) > t$. This formula together with the electrostatic force contributes to achieve the aesthetic criterion 3 (see Figure 9).

For each iteration and pair of nodes, the force that moves the node is equal to the sum of the aforementioned forces.

Rejected Force.

After a certain number of iterations, an additional force is used to avoid the overlapping of the intermediate nodes with the target nodes.

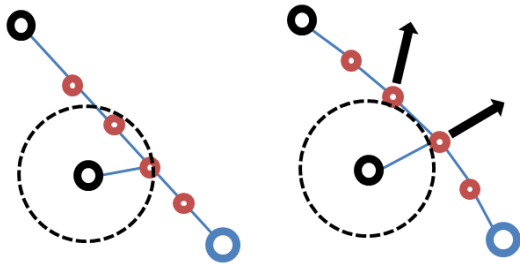


Figure 7: Rejected force applied to the intermediate nodes of a flow.

$$F_r(p_i, q_i) = \frac{-1}{\|p_i - q_i\|} (\widehat{p_i - q_i}) \quad (4)$$

This force is in magnitude equivalent to the electrostatic force, but it has opposite direction (see Figure 7). Moreover, only intermediate nodes near to a certain distance are affected by this force. This formula satisfies the aesthetic criterion 4 "The target nodes are not overlapped with flows".

In this case, for each iteration and for each couple of nodes, the following formula is calculated:

$$F_f(p_i) = \sum_{s \in I_{p_i}} F_e(p_i, s) + \sum_{t \in T} F_r(p_i, t) + F_s(p_i) \quad (5)$$

I_{p_i} is the set of intermediate nodes that interact with p_i , i.e. they have the same index and the same father. T is the set containing all the target nodes. In order to balance all the forces we introduce one constant for each force. Finally, these are the two final formula used by the algorithm, one with the rejected force and the other without it:

$$F_f(p_i) = k_e \times \sum_{s \in I_{p_i}} F_e(p_i, s) + k_r \times \sum_{t \in T} F_r(p_i, t) + k_s \times F_s(p_i) \quad (6)$$

$$F_f(p_i) = k_e \times \sum_{s \in I_{p_i}} F_e(p_i, s) + k_s \times F_s(p_i) \quad (7)$$

k_e and k_s are constants defined by the user. k_r can be assigned as $2k_e$ since the reject force has to be greater than the electrostatic force. The combination

of rejected force, electrostatic force and stressed force creates 'smooth' flows despite its composition by straight segments.



Figure 9: Flow maps of immigration from Colorado. On top without the stress force, in the middle with the simple electrostatic force using formula (3), on bottom the stress force using a threshold, note in the last case as the main flows are represented as smooth.

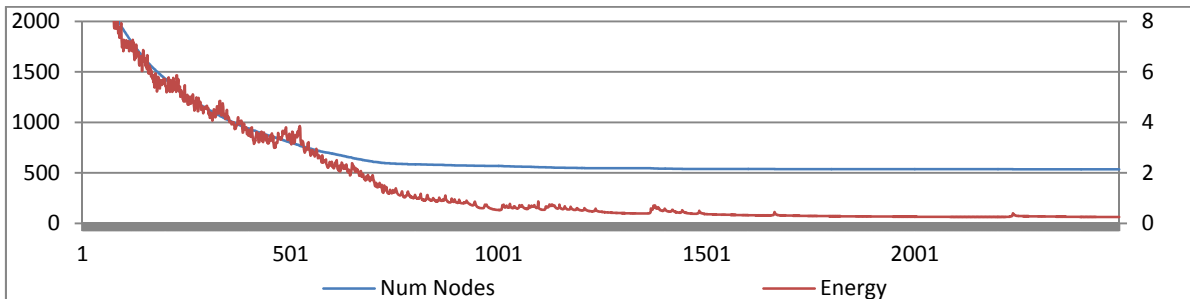


Figure 8: Chart showing the number of nodes (blue line, left axis) and the energy (red line, right axis) at every iteration (X axis).

Hence it fulfills the aesthetic criterion 2.

In this phase the displacements of all nodes are calculated (accordingly with the forces) and then these displacements are applied to the nodes.

3.3 Rendering Phase

To improve the visual appealing, we tested a few rendering settings. First we represented flow lines using parametric curves, in particular the natural cubic spline (Spcah, 1974), using the intermediate nodes with more than one child as control points. However the result achieved was not satisfactory. Therefore this phase is quite simple; a width line is created from each node to its child. The thickness of the line identifies the magnitude of its child. Hence, this algorithm does not require the use of parametric curves.

It is the sum of the forces and the high number of nodes that generates a visual appealing result. The size of the target nodes is proportional to the flow magnitude.

4 RESULTS, DATASETS & IMPLEMENTATION

Our program computes an optimal flow tree for most maps in a couple of minutes on a system with 1.64 GHz and 2 GB of RAM. Flow trees are automatically generated by our program without requiring any user intervention.

The algorithm was initially developed using a 2D Java canvas, but was later integrated in NASA WorldWind (NASA, 2013).

The complexity of the algorithm is $O(i \times n^2)$, where i is the number of iterations and n is number of nodes. To benchmark our algorithm, we decided to use a few datasets. The results presented in this section corresponds the dataset describing the migration from California. We decided to pick this dataset because we want to compare our results to the ones of previous algorithms (see Figure 2). Accordingly with the initial number of 4176 intermediate nodes, it takes 227 seconds to have a good result.

The chart in Figure 8 describes at each interaction, the trend of the number of intermediate nodes and the energy, that is, the sum of all the forces, see formula (7). It is possible to notice that the number of nodes is stable after 1200 iterations meanwhile the energy is stable after 2300 iterations. When the energy stability is reached, the algorithm

applies the rejected force, see formula (6). We apply the rejected force at this stage because during the first iterations this force can avoid the merging of flows. This result and the time needed depend mostly on two different factors: the number of initial intermediate nodes and the constants used for the different forces.

Another dataset describes the relation between Fondazione Graphitech and its partners. The magnitude defines the number of shared projects. Moreover, the exports of Italian goods in all over the World were also represented (see Figure 1).

5 CONCLUSION AND FUTURE WORK

We have presented a force-directed algorithm for the automatic generation of flow maps: by applying a specific set of forces to the set of flow lines, we can easily satisfy all criteria that characterize a well-drawn flow map.

Additionally, we provide an evaluation of the system that also shows that results are clearly much better than in previous approaches.

There are a few aspects in this work that have still margin for improvements. The performance can be increased using the GPU (Frishman and Tal, 2007). A procedure to quantify the quality of the visualization comparing the result with other techniques should be performed. For example measuring the amount of pixel-based overdraw and empty space.

During the rendering phase an exhaustive study has to be done to decide the more appropriate parametric curve with high visual appealing.

At the moment the user has to define manually the force factors; a possible improvement could be the possibility to auto assign all the parameters accordingly to the scale of the dataset.

Nowadays the users value the possibility to interact with the maps. The force-directed algorithm can be a powerful tool within such interactive process. For example moving an intermediate node, the others automatically are moved improving the visual quality.

Additionally, this algorithm can be extended taking into account more than one origin. The key aspect will be the reduction of intersections between flows.

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