

HIDE: Hybrid Symmetric Key Algorithm for Integrity Check, Dynamic Key Generation and Encryption

Jayagopal Narayanaswamy, Raghav V. Sampangi and Srinivas Sampalli
Faculty of Computer Science, Dalhousie University, Halifax, Nova Scotia, Canada

Keywords: Dynamic Key Generation, Chained Key Generation, Encryption, Integrity Check.

Abstract: This paper proposes a hybrid encryption technique that generates a key dynamically, along with integrity check parameters. Our approach generates the key stream using a chained approach, beginning with an initial pre-shared key. Subsequent keys are derived using logical operations on intermediate cipher texts and intermediate keys generated in each stage. This is an improvement over chaining techniques, which use a cipher text to derive successive keys. We validate our algorithm by proof-of-concept implementation and security analysis.

1 INTRODUCTION

Block ciphers split a message into blocks of identical sizes, which are encrypted with key blocks. A key for a block cipher remains the same or is derived from an initial key using functions such as the Feistel function (Knudsen, 1994). In contrast, stream ciphers encrypt a stream of bits from a message with a key bit-by-bit using the logical exclusive-OR (XOR) operation. The key stream is generated by a PRNG through a fixed size input called the seed.

In stream ciphers, the encryption operation is simple but vulnerable to attacks such as distinguishing and key recovery attacks (Hell et al., 2009). Stream ciphers generate a continuous stream of key bits using a Linear Feedback Shift Register (LFSR) that require an n -bit seed value. A drawback with a LFSR is that an n -bit pattern may repeat in the key stream before completing all 2^n possible patterns (Zeng et al., 1991). To avoid this issue a Non-Linear Feedback Shift Register (NLFSR) has been proposed, but there are no generic designs for a NLFSR (Dubrova et al., 2008) (Bardis et al., 2004). In addition, in both block and stream ciphers, the key is derived from an initial key or a seed, therefore, knowledge of the initial key reveals information about the original message. To overcome this issue, we use multiple keys that are generated dynamically.

Recently, an algorithm called Hummingbird was proposed as a hybrid of the stream and block cipher approaches that was designed for resource-constrained devices, such as, Radio Frequency Identification (RFID) systems (Engels et al., 2010). Hum-

mingbird follows the traditional encryption process of block ciphers such as substitution and looping; while a key is derived through the stream cipher principle. The key generation and encryption process in Hummingbird is computationally expensive. However, we compare our algorithm with Hummingbird as an example. We don't inherit any property of the Humming algorithm in our approach.

After the encryption process, any encrypted message is hashed to offer integrity. Hashing is a one-way mathematical function, which converts the variably sized messages into a fixed size output called a Message Digest (MD). The receiver verifies the integrity of received message by hashing it and then compares it with the received MD.

In this paper, we propose a hybrid encryption algorithm that requires a simple (Exclusive-OR) XOR operation. We use the stream cipher approach to derive a key, while the block cipher approach is adopted in the encryption process. In our approach, a key stream is generated from a previous key block and an intermediate cipher text block, which is encrypted with a message block-by-block using the XOR operation. In addition to encryption, a fixed-size final key in each round is used as the MD that provides the integrity to the message. Since we use only the XOR operation throughout the encryption process and provide the integrity check parameter without using any external hashing algorithm, we expect that our approach will reduce the computational complexity as well as increase performance. We have implemented the algorithm and we show that our algorithm can withstand such potential attacks as differential,

known cipher text, known plain text, distinguishing and key recovery attacks. We also conducted standard security analysis tests (such as, entropy analysis, periodicity check, frequency test, poker test, run test, serial test etc.) and have included the results in the security analysis section.

2 BACKGROUND AND LITERATURE REVIEW

In stream ciphers, an LFSR plays an important role in the generation of a key stream. One possible implementation of a Pseudo Random Number Generator (PRNG) is using an LFSR. An LFSR can be constructed by the consecutive assembly of shift registers. An n -bit LFSR produces a continuous bit stream from an n -bit seed. One of the drawbacks of an LFSR is that it will form a cycle on or before reaching 2^n patterns for a given n -bit seed (i.e. the same sequence of bits will be generated on or before reaching 2^n patterns) (Zeng et al., 1991).

To overcome this drawback, an NLFSR was introduced (Dubrova et al., 2008) (Bardis et al., 2004). NLFSRs avoid the linearity problem and extend the cyclic period. A common way of constructing an NLFSR is to use more than one LFSR connected through logic gates. Logic gates are used to choose the current LFSR output from a list of connected LFSRs. Thus, linearity can be broken easily and the cyclic period will be extended. Even though NLFSR acts as an alternative to LFSR, there is no NLFSR that guarantees a long cyclic period. Furthermore, a general technique for constructing NLFSRs has been an open problem (Lv et al., 2012) (Rachwalik et al., 2012) (Hu and Gong, 2011) (Mandal and Gong, 2012).

Block ciphers offer better security over stream ciphers, however the latter are computationally more efficient. Hybrid approaches that combine the benefits of both have been the focus of recent research. Trivium (De Cannière, 2006) is the stream cipher that uses the block cipher principle for the key generation process. The Hummingbird algorithm encrypts the message block-by-block and uses internal registers to update the key stream. It has a small block size of 64-bits, which is further divided into four smaller blocks of 16-bits each that are equivalent to the number of bits in the internal state register. Initially, the internal state register is loaded with random bits. Furthermore, the LFSR and the cipher text are used to regenerate the internal state register bits. It uses a 256-bit key that is split into four 64-bit keys to encrypt the message. Two more secret keys, namely, $K5$ and

$K6$, are introduced and these are derived from the initial four keys. The Hummingbird algorithm has four rounds of substitution and permutation and uses four keys on each round to encrypt the message. In the fifth round, the output is encrypted with key $K5$ followed by a final substitution and then encryption with key $K6$. This approach is more suitable for resource-constrained devices such as RFID systems. It is resistant to differential as well as linear attacks (Engels et al., 2010). However, we use this algorithm as an example of a hybrid approach and our algorithm does not inherit any properties used in Hummingbird.

Kutuboddin and Krupa (Jinabade and Rasane, 2013) have extended the research on the Hummingbird algorithm to increase its computational efficiency. They propose an approach, which replaces the traditional Hummingbird modulo operation by introducing Field Programmable Gate Array (FPGA). As the proposed approach uses XOR operation, computational efficiency is better compared to the original version.

Integrity is an important component of network security that prevents an anonymous entity from data manipulation. The Integrity of the message can be achieved through hashing algorithms, such as, MD5, SHA1, SHA2, SHA3, etc. Traditional hashing is a mathematical one-way function, which encrypts the message but decryption is not possible. It converts any variable size message to a fixed size output called the Message Digest (MD). A collision between two MD is possible after hashing 2^n messages. However, as the encryption (or decryption) and hashing are disjoint operations, the use of hashing will lead to computational overhead on the system (Kumar et al., 2011) (Hu and Wang, 2009) (Li et al., 2012).

In recent years, a hybrid approach called Authenticated Encryption (AE) has provided integrity as well as authenticity for short messages (Jeddi et al., 2013). However, it requires a Message Authentication Code (MAC), which is derived from the hashing technique. This creates further computational overhead. For example, Signcryption was proposed for the asymmetric key algorithm (Zheng, 1997)(Wang et al., 2013), which replaced a traditional encrypt and sign practice by integrating message signatures (similar to MD) as a part of an encryption process. Zheng et al.'s (Zheng, 1997) work on the signcryption scheme is based on the theory that a combined computational cost of a signature (using hashing) and encryption will be less than their individual costs. In signcryption, part of the key is generated using a hashing algorithm to provide the integrity, and the same key is used to encrypt (using an asymmetric key algorithm) the message to offer confidentiality, integrity and authentica-

tion (Zheng, 1997).

Our literature review concludes that existing algorithms are weak in generating a key or the encryption process leads to computational overhead. To fill a gap between computational overhead and to support a strong dynamic key generation process to encrypt the message, we propose a new approach that uses a simple XOR operation for better computation, and Intermediate Cipher Text (ICT) to generate the key dynamically.

3 PROPOSED APPROACH

3.1 Overview

Table 1 contains a list of notations that are used in the proposed approach.

In block ciphers, common modes of encryption such as Cipher Block Chaining (CBC), Propagating Cipher Block Chaining (PCBC) and Cipher Feedback (CFB) use the cipher text as part of the encryption. This may lead to some information about the original message being revealed to an attacker. To avoid this issue, our approach uses an Intermediate Cipher Text (ICT) to generate a dynamic key. Since the dynamic key is used on each block of the message encryption, the cipher text bits are completely random.

Table 1: Notations.

K_a	Keys to encrypt in Round One
K_b	Keys to encrypt in Round Two
n	Last block
l	Current block
i_x	Prediction bit
M	Message to be encrypted
$m_1, m_2 \dots m_l$	Message blocks
XOR	Logical bit-by-bit XOR operation
$\delta()$	Derivation function
$\rho()$	Prediction function
$\gamma()$	Inversion function
M_x	Message block x
K_{ax}	Round One key block x
K_{bx}	Round Two key block x
ICT_x	Intermediate Cipher Text block x
CT_x	Cipher Text block x
$C_{\alpha, \beta}$	Chunk β of component α (α can be $M_x, ICT_x, K_{ax}, K_{bx}, CT_x$)
KG1	Key generation for Round One
KG2	Key generation for Round Two

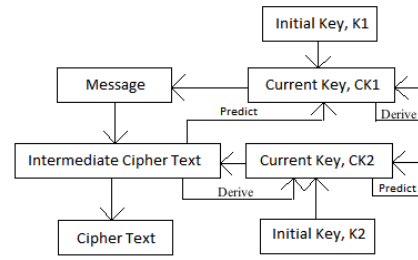


Figure 1: Overview of the encryption process.

A hybrid encryption algorithm is proposed, which includes both stream cipher and block cipher features. This approach consists of two parts: a key generation part (based on stream ciphers) and two rounds of encryption (or decryption) using a basic XOR operation (employing message blocks as in block ciphers). The message is encrypted block-by-block with different (block) keys on each message block. The size of one block in the message is 128-bits in sixteen 8-bit chunks.

Figure 1 illustrates the overall encryption process. This approach requires two initial keys that encrypt the first block in each round. Except for the initial key, every successive key (to encrypt each block) is derived from the bits in the current key block and an ICT block. In general, blocks of message are encrypted using blocks of keys generated using the stream cipher scheme. Key generation and encryption follows a chained approach, with the current (block) keys and ICTs used to generate (block) keys for the encryption of the subsequent message blocks. The encryption and key generation processes are described later in this section.

3.2 Assumptions

It is assumed that the initial key used to encrypt the first block of the message in each round is generated by the sender and securely transmitted to the receiver. The key used to encrypt the message is randomly chosen from a set of strong keys.

3.3 Design

3.3.1 Key Generation

Since it is assumed that the key is known only to the sender and the receiver, we generate successive keys from the previous keys with the help of an intermediate cipher text (ICT). The term “intermediate cipher text” represents an encrypted message after Round 1, whereas the actual cipher text is generated after Round 2 (as will be discussed in detail in the encryp-

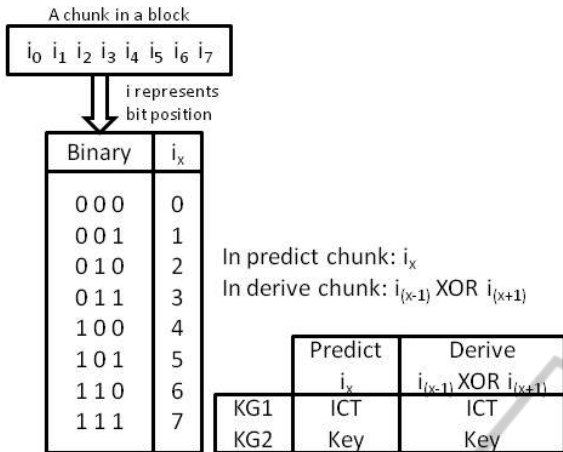


Figure 2: Key Generation Process Overview.

section). The key generation process in this approach is different for both rounds. Each round uses a different initial key (K_{a0} & K_{b0}) to encrypt the first block in the message. To encrypt each successive block in the message, we introduce a combination of *prediction* and *derivation* techniques to generate successive keys ($K_{a1...n}$ for Round 1 and $K_{b1...n}$ for Round 2) from a previous key block and an ICT block.

In *prediction*, three bits in every chunk per block are used to choose the i_x^{th} bit within the respective chunk. The value of i_x is based on the three binary bits chosen in the chunk. The value of i_x is a decimal representation of the 3-bit binary. In a chunk, 3-bits are selected in a clockwise direction starting from the most significant bit (MSB) in the chunk to the least significant bit (LSB) to choose one out of eight possible values for i_x , which is repeated 8 times starting from the MSB to the LSB. The value of i represents the position of the binary bit in the chunk. Each value of i_x predicts a new bit that is used to generate chunks per block in every successive key. As shown in Figure 2, three bits are rotated in a clockwise direction to yield one out of eight possible values for i_x .

In *derivation*, the successive key is generated by XORing the $(i_x - 1)^{\text{th}}$ position of the binary bit with the $(i_x + 1)^{\text{th}}$ position of the binary bit to form a new bit. The same process is repeated in every chunk per block to generate a whole key for encrypting the successive block in the message. The whole process of key generation offers a randomness in the key stream without forming a cycle.

The process for applying *prediction* and *derivation* in combination to generate a key is as follows. In Round 1, the key is derived from key K_a series based on the ICT block's prediction. In contrast, Round 2's successive keys are derived from the ICT block based on the key block's prediction.

To generate a key, the *prediction* and *derivation* combination is applied alternatively on an ICT block and a key block over two rounds.

3.3.2 Encryption (or Decryption)

Initially, in Round 1, four out of eight bits (say 0, 2, 3, 5) in every chunk in a key K_{a0} are inverted. The first chunk in the initial key K_{a0} is XORed with the first chunk in the first block of the message, which produces the first chunk in ICT_0 (which represents the ICT first block). Again, four out of eight bits (say 2, 4, 5, 7) in every chunk in ICT_0 are inverted (this process is common throughout the ICT blocks), which is followed by successive key generation operations (already discussed in detail earlier). The output of the key generation operation will be the first chunk in the key K_{a1} to encrypt the first chunk of the next message block. The first chunk of key K_{a1} is XORed with the second chunk of key K_{a0} that will be used to encrypt the second chunk in the first block of the message. In conclusion, every first chunk in each block of the message is directly XORed with the first chunk of the respective key, whereas the successive chunks in each block in the message are encrypted with the XORed output of the successive chunk in the present key and the currently derived chunk for the next key. For example, the 4^{th} chunk in message $C_{M2,4}$ will be encrypted with the XORed output of the 4^{th} chunk in the current key $C_{K_{a2,4}}$ and the 3^{rd} chunk in the next key $C_{K_{a3,3}}$.

Round 2 follows the same procedure as Round 1, except that the message and initial key K_{a0} will be replaced by the ICT and key K_{b0} respectively. In addition, the inversion operation is applied on bits 1, 4, 6, 7 and 0, 1, 3, 6 for the key and the ICT respectively. The output of Round 2 is the final cipher text. Every block in a message is XORed with the key block to generate a cipher text. Figure 3 illustrates the encryption process.

A decryption process is symmetric with encryption in reverse manner.

3.3.3 Integrity Check

The final keys (say $K_{a(n+1)}$ and $K_{b(n+1)}$) in each round of encryption will act as the integrity check parameter. In this approach, the 128-bit key $K_{a(n+1)}$ will be concatenated with the 128-bit key $K_{b(n+1)}$ to form a 256-bit message digest, which is used for integrity check.

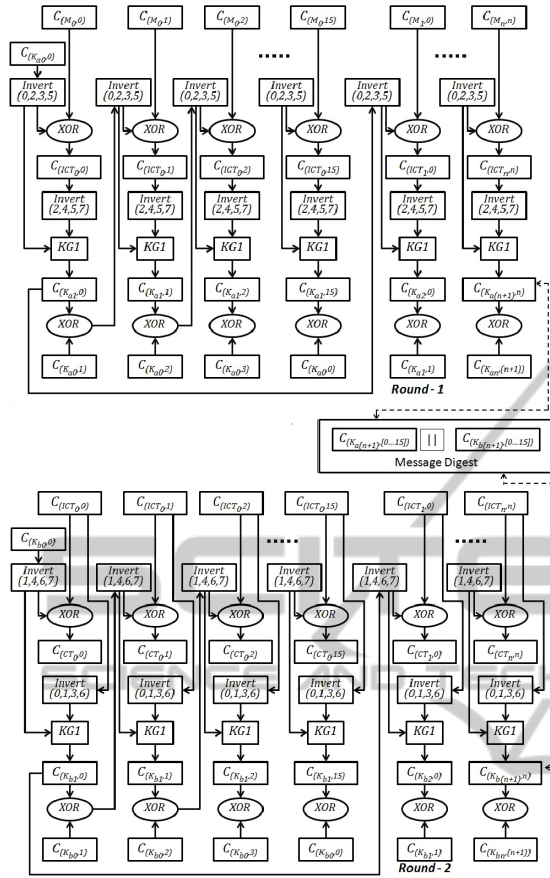


Figure 3: Encryption Process.

3.4 Example

3.4.1 Encryption and Key Generation

In this section the two rounds of the encryption process will be explained with a generic example. Each round used different initial key (K_{a0} and K_{b0}). Further, the message (M) is segmented into different blocks. The first block of the plain text message is encrypted with a key combination of K_{a0} (the present key) and K_{a1} (the currently derived key) at Round 1 that generates the first ICT block.

Let the message blocks be:

$$M = M_0, M_1, M_2, M_3, \dots, M_n \quad (1)$$

Round – 1 : Encrypting the first block of the message: The message block will be encrypted chunk by chunk. Before the XOR operation, bit positions 0, 2, 3, 5 will be inverted in the first chunk of an initial key. Then, the first chunk in the first block of the message will be XORed with the initial key's first chunk (which is after the inversion).

$$C'_{K_{a0},0} = \gamma(C_{K_{a0},0}) \quad (2)$$

$$C_{ICT0,0} = E_{C'_{K_{a0},0}}(C_{M0,0}) \quad (3)$$

Once again, bits 2, 4, 5 and 7 in the ICT will be inverted before the key generation process and the output will be used to generate a key to encrypt the next chunk as well as the block.

$$C'_{ICT0,0} = \gamma(C_{ICT0,0}) \quad (4)$$

Each bit of the key, $C_{K_{a1},0}$, is generated from the previous key, $C'_{K_{a0},0}$, whereas, n bits in $C'_{ICT0,0}$ chooses a bit in $C'_{K_{a0},0}$ to produce the new key, $C_{K_{a1},0}$. In general, a current key, K_{al} , is generated from the prediction of a previous ICT block, ICT_{l-1} , and the derivation of the previous key, $K_{a(l-1)}$, in Round 1. Every first chunk in the message block is directly XORed with the inversion of the first chunk in the key, whereas the 2^{nd} to 15^{th} chunk in the message will be encrypted with an inversion of the XORed value of the next chunk key with next block's chunk's key.

Key Generation: for initial chunks only

$$C_{K_{al},0} = \rho(C'_{ICT_{l-1},0}[\delta(C'_{K_{a(l-1)},0})]) \quad (5)$$

For the 2^{nd} to 15^{th} chunks

$$C_{K_{al},1} = [C_{K_{al},1}]XOR[C_{K_{a(l+1)},0}] \quad (6)$$

...

$$C_{K_{al},15} = [C_{K_{al},15}]XOR[C_{K_{a(l+1)},14}] \quad (7)$$

The generic equation for encryption in Round 1 is as follows:

$$ICT_0 = E_{K_{a0}}(M_0), ICT_1 = E_{K_{a1}}(M_1), \dots, ICT_n = E_{K_{an}}(M_n) \quad (8)$$

Figure 4 is an example diagram for the encryption and integrity check processes. For this example, it is assume that the message contains only two blocks and that each one has only two chunks (the typical block size used in this example is 16-bits, which is composed of two 8-bit chunks). The initial key, K_{a0} (10011010 11011011), is encrypted with the message (00101101 00110011 10111011 01000011) which produces, in Round 1, the ICT (00000011 11010010 10000001 01001111). Then, in Round 1, the ICT block is used to predict a derivation bit in a key block, whereas in Round 2 a key block predicts the derivation bit in the ICT block. Further, in Round 2, the cipher text (11110011 01000111 00111000 10111101) will be obtained by encrypting the ICT with the key block K_{b0} (10111011 00101100). Finally, the concatenation of the last derived keys, K_{a2} and K_{b2} , will act as the message digest (01000111 11100010 10111001 11001111).

Round – 2 : In this Round, the encryption process is the same as in Round 1. But in the key generation

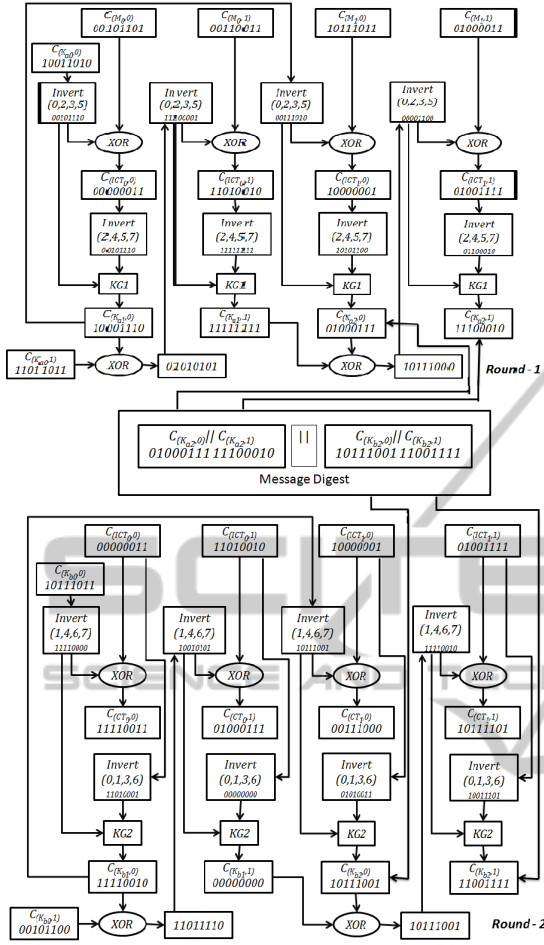


Figure 4: An example of the Encryption and Integrity Check Process.

process, the prediction and derivation functions are interchanged between an ICT block and a key. An initial key (K_{b_0}) encrypts the first block of the ICT block. To encrypt subsequent blocks, a current key (K_{bl}) is generated from the prediction of a previous Key block ($K_{b(l-1)}$) and the derivation of a previous ICT block (ICT_{l-1}). In Round 2, the bits chosen to invert are different from Round 1. For a key bits 1, 4, 6 and 7 are inverted before encryption, whereas bits 0, 1, 3 and 6 are inverted for the ICT chunks (after the encryption, but before the key generation process).

Encrypting first block of the message:

$$C'_{K_{b_0},0} = \gamma(C_{K_{b_0},0}) \quad (9)$$

$$C_{CT_0,0} = E_{C'_{K_{b_0},0}}(C_{ICT_0,0}) \quad (10)$$

$$C'_{ICT_0,0} = \gamma(C_{ICT_0,0}) \quad (11)$$

Key Generation: for initial chunks only

$$C_{K_{bl},0} = \rho(C'_{K_{b(l-1)},0}[\delta(C'_{ICT_{l-1},0})]) \quad (12)$$

For the 2nd to 15th chunks

$$C_{K_{bl},1} = [C_{K_{bl},1}]XOR[C_{K_{b(l+1)l},0}] \quad (13)$$

...

$$C_{K_{bl},15} = [C_{K_{bl},15}]XOR[C_{K_{b(l+1)l},14}] \quad (14)$$

The generic equation for encryption in Round 2 is as follows:

$$CT_0 = E_{K_{b_0}}(ICT_0), CT_1 = E_{K_{b_1}}(ICT_1), \dots, CT_n = E_{K_{b_n}}(ICT_n) \quad (15)$$

3.4.2 Message Digest

The final keys generated from each round will be concatenated to yield the 256-bit Message Digest.

$$MD = K_{a(n+1)} || K_{b(n+1)} \quad (16)$$

4 ANALYSIS AND DISCUSSION

This proposal has successive key derivation function, which uses both a key as well as a message to choose the next key dynamically. Every message block is encrypted with a different key chosen from the 2^{128} possible combination keys for a given 128-bits in the key space. As successive keys are generated based on a message and a key in any one process of either prediction or derivation, the generated key is completely dynamic. Thus, it guarantees that there would be no regular cycle in a key bit stream. The algorithm is a very simple, which uses XOR operations to generate keys as well as encrypt a message. It can withstand a variety of attacks (chosen-plain text attack, chosen-cipher attack, differential attack, linear attack and distinguishing attack) (Hell et al., 2009) (Zulkifli, 2008) (Northwood, online).

4.1 Cryptanalysis

A short description of possible attacks and their analyses are explained in this section. Chosen-plain text and chosen-cipher text attacks are targeted on symmetric key encryption schemes. Both attacks work on the principle of choosing a piece of information for retrieving an original message for a given cipher text or to reach a cipher text for a given plain text. In this approach, since the key is changed for every block encrypted, chosen-plain text and chosen-cipher text attacks are hard to launch. A differential attack compares an input value with an output value to obtain a possible key. Since this proposal relies on both a key and an ICT, and the key is chosen from a strong key, a differential attack is difficult to implement. In

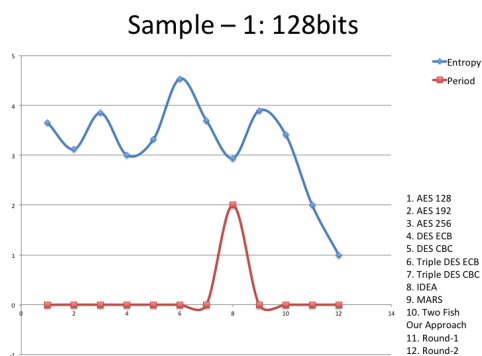


Figure 5: Sample 1.

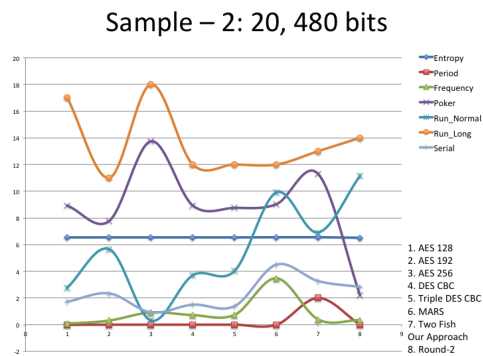


Figure 6: Sample 2.

linear cryptanalysis, an equation is formed from plain text and cipher text, which is equated to possible key bits to reveal some information about the message. To avoid this issue, the key generation process is switched between prediction and derivation techniques. A distinguishing attack focuses on stream ciphers that compare a given sequence of values to check the randomness. A dynamically generated key ensures that there will be no relation between the current and previous keys, so launching a distinguishing attack is difficult.

4.2 Security

In addition to the attacks mentioned above, tests were done using a cryptanalysis tool called Cryptool (url, 2014). As mentioned earlier, though the key generation part is in the nature of a stream cipher, the actual encryption process adopts a block cipher approach. The security analysis of the proposed approach was compared with existing block cipher algorithms such as AES, DES, IDEA, MARS, Twofish etc. (the implementation of which is available as a built in option in Cryptool). Graphs are presented below with five different scenarios, which are based on the cryptanalysis tests (entropy test, periodicity test, frequency test, poker test, run test and serial test) (Wu et al., 2011) (Biebighauser, 2000) (Soto, online). The results from this analysis show that our algorithm passes the above mentioned tests and that the security level is no lower than the pool of example algorithms used. Figures 5, 6, 7 display the sample graphs that were taken from the security analysis with different sample inputs.

The proposed algorithm provides non-linearity in the key cycle, which is achieved by using an intermediate cipher text block. However, using cipher text to encrypt a message is vulnerable, so two rounds of encryption were designed with an ICT as the second round key generator instead of the cipher text. Applying key generation from a key for the first round and an ICT for the second round ensures a dynamic property in successive keys.

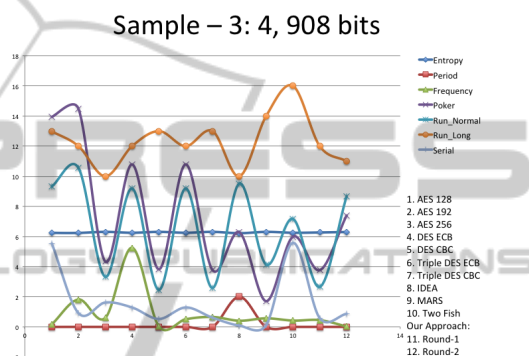


Figure 7: Sample 3.

4.3 Performance

The performance of our approach has yet to be tested, but it is expected to be computationally efficient, which is our contention based on the following claims. Initially, the computation is required only for the encryption part, but the integrity check does not require any additional computation. As the key generation process is a bitwise operation, it is the only difficulty in this approach. The number of operations required per bit is only fourteen, which includes encryption (or decryption), an integrity check and key generation. We identified from our research that an increase in performance would require an increase in the hardware component, which may not be appropriate depending on the requirements of the application. However, we are currently working on the key generation part.

5 CONCLUSION AND FUTURE WORK

In this paper, a hybrid encryption algorithm is proposed with increased security. In our approach, blocks of message are encrypted using the blocks of

keys generated by a stream cipher scheme. Key generation and encryption follows a chained approach in which current (block) keys and ICTs are used to generate (block) keys for the encryption of subsequent message blocks. This approach is feasible for encrypting a message using a hybrid stream cipher and block cipher approach. The novelty of our approach is its dynamic key generation from an intermediate cipher text with an the integration of an integrity check. We are currently extending our research to optimize the computational efficiency of our approach.

REFERENCES

- (2014). The cryptool portal @ONLINE. <http://www.cryptool.org/en>.
- Bardis, N., Markovskyy, A., and Andrikou, D. (2004). Method for designing pseudorandom binary sequences generators on nonlinear feedback shift register(nfsr). *WSEAS Transactions on Communications*, 3(2):758–763.
- Biebighauser, D. (2000). Testing random number generators @ONLINE. <http://www-users.math.umn.edu/~garrett/students/reu/pRNGs.pdf>.
- De Cannière, C. (2006). Trivium: A stream cipher construction inspired by block cipher design principles. In *Information Security*, pages 171–186. Springer.
- Dubrova, E., Teslenko, M., and Tenhunen, H. (2008). On analysis and synthesis of (n, k) -non-linear feedback shift registers. In *Design, Automation and Test in Europe, 2008. DATE'08*, pages 1286–1291. IEEE.
- Engels, D., Fan, X., Gong, G., Hu, H., and Smith, E. M. (2010). Hummingbird: ultra-lightweight cryptography for resource-constrained devices. In *Financial Cryptography and Data Security*, pages 3–18. Springer.
- Hell, M., Johansson, T., and Brynielsson, L. (2009). An overview of distinguishing attacks on stream ciphers. *Cryptography and Communications*, 1(1):71–94.
- Hu, H. and Gong, G. (2011). Periods on two kinds of nonlinear feedback shift registers with time varying feedback functions. *International Journal of Foundations of Computer Science*, 22(06):1317–1329.
- Hu, M. and Wang, Y. (2009). The collision rate tests of two known message digest algorithms. In *Computational Intelligence and Security, 2009. CIS'09. International Conference on*, volume 2, pages 319–323. IEEE.
- Jeddi, Z., Amini, E., and Bayoumi, M. (2013). A novel authenticated encryption algorithm for rfid systems. In *Digital System Design (DSD), 2013 Euromicro Conference on*, pages 658–661. IEEE.
- Jinabade, K. and Rasane, K. (2013). Efficient implementation of hummingbird cryptographic algorithm on a reconfigurable platform. In *International Journal of Engineering Research and Technology*, volume 2. ESRSA Publications.
- Knudsen, L. R. (1994). Practically secure feistel ciphers. In *Fast Software Encryption*, pages 211–221. Springer.
- Kumar, Y., Munjal, R., and Sharma, H. (2011). Comparison of symmetric and asymmetric cryptography with existing vulnerabilities and countermeasures. *International Journal of Computer Science and Management Studies*, 11(03).
- Li, X., Zhang, W., Wang, X., and Li, M. (2012). Novel convertible authenticated encryption schemes without using hash functions. In *Computer Science and Automation Engineering (CSAE), 2012 IEEE International Conference on*, volume 1, pages 504–508. IEEE.
- Lv, H., Xie, J.-X., Fang, J.-C., and Qi, P. (2012). Generating of a nonlinear pseudorandom sequence using linear feedback shift register. In *ICT Convergence (ICTC), 2012 International Conference on*, pages 432–435. IEEE.
- Mandal, K. and Gong, G. (2012). Probabilistic generation of good span n sequences from nonlinear feedback shift registers. *University of Waterloo*.
- Northwood, C. Cryptography, attacks and countermeasures @ONLINE. <http://www.pling.org.uk/cs/cry.html>.
- Rachwalik, T., Szmidi, J., Wicik, R., and Zablocki, J. (2012). Generation of nonlinear feedback shift registers with special-purpose hardware. In *Communications and Information Systems Conference (MCC), 2012 Military*, pages 1–4. IEEE.
- Soto, J. Statistical testing of random number generators @ONLINE. <http://infosec.pku.edu.cn/~tly/oldversion/nist-nissc-1999/papers/p24.pdf>.
- Wang, Y., Manulis, M., Au, M. H., and Susilo, W. (2013). Relations among privacy notions for signcryption and key invisible sign-then-encrypt. In *Information Security and Privacy*, pages 187–202. Springer.
- Wu, Y., Noonan, J. P., and Agaian, S. (2011). A novel information entropy based randomness test for image encryption. In *Systems, Man, and Cybernetics (SMC), 2011 IEEE International Conference on*, pages 2676–2680. IEEE.
- Zeng, K., Yang, C.-H., Wei, D.-Y., and Rao, T. (1991). Pseudorandom bit generators in stream-cipher cryptography. *Computer*, 24(2):8–17.
- Zheng, Y. (1997). Digital signcryption or how to achieve cost (signature & encryption) \leq cost (signature) + cost(encryption). In *Advances in Cryptology-CRYPTO'97*, pages 165–179. Springer.
- Zulkifli, M. Z. W. M. (2008). Attack on cryptography @ONLINE. <https://idazuwaika.files.wordpress.com/2008/06/attack-on-cryptography.pdf>.