

Dy-COPECA: A Dynamic Version of MC/DC Analyzer for C Program

Sangharatna Godbole¹ ^a and Arpita Dutta² ^b

¹Department of CSE, NIT Warangal, Telangana, India

²Department of CSE, IIT Kharagpur, West Bengal, India

Keywords: MC/DC, Test Cases, Software Testing, Static MC/DC Analysis, Dynamic MC/DC Analysis.

Abstract: RTCA/DO-178B&C standards mandate Modified Condition / Decision Coverage (MC/DC) criterion for level-A category software. In critical safety system applications such as Aircraft or Metro Rail controller systems, etc., testing engineers have to produce the MC/DC report. There are several MC/DC analyzers, which are either automated or partially-automated available. Some of the existing analyzers do not consider the dependencies of Predicates/Decisions on each other. These analyzers process each predicate individually based on MC/DC criterion. They use test cases to identify the total number of atomic conditions present in a decision which influence the output of whole decision. In this paper, we overcome the limitations of some of the existing techniques. We propose an approach, which execute the whole program along with unit test cases at run time to compute MC/DC score. This dynamic mechanism solves the dependency relation between the variables appearing at different predicates and their branch statements in a single run. We have developed Dynamic COverage PErcentage CAculator (Dy-COPECA) using C and Java language to process C-programs. We have improved the MC/DC by 42.88% through dynamic MC/DC analysis as compared to static analysis for the example C-program.

1 INTRODUCTION


Software Testing is an important phase of Software Development Life Cycle (SDLC). Manual software testing accounts for 50-80% of total software development cost (Myers et al., 2011; Beizer, 2003; Chauhan, 2010; Mall, 2018). Manually created test cases and computing code coverage are expensive (Gao et al., 2005), error-prone, and generally not exhaustive (Majumdar and Sen, 2007). Therefore, automated software testing techniques have been discovered (Bird and Munoz, 1983; Csallner et al., 2008; Gupta et al., 1998; DeMillo and Offutt, 1993).


White-box testing is one of the types of software testing techniques (Ammann et al., 2003). White-box testing deals with structural testing, where testers have the knowledge and resources in terms of source code. There are several code coverage criteria (Grindal et al., 2005; Rajan et al., 2008) are available. Out of which MC/DC is the second strongest criterion which requires minimum “n+1” number of test case and maximum “2n” number of test cases, where “n” is the total number of atomic conditions

present in the predicates of a program.

MC/DC was proposed a few decades ago (Chilenski and Miller, 1994; Jones and Harold, 2003). Many researchers across the globe are considering it as an important technique. There are also several commercialized organizations which reports MC/DC for critical industrial software applications. Different researchers have implemented MC/DC analyzer, but those have several limitations. The main disadvantage of measuring MC/DC by the existing techniques is their static mechanism. As a result, those analyzers are failed to measure the correct MC/DC for a program. We have also observed that several existing analysers work with high manual intervention which means, they are semi-automated.

In this paper, we report this issue of static mechanism to measure MC/DC. The existing analyzer applies MC/DC mechanism on all the predicates present in a program individually. These existing works do not care for the functional effect of variables, atomic conditions, and predicates which appear one after other in a program. It means one predicate appears in a program before another predicate without taking care of updated values for variables appeared in sec-

^a  <https://orcid.org/0000-0002-6169-6334>

^b  <https://orcid.org/0000-0001-7887-3264>

ond predicate. When we compute MC/DC for these predicates with static mechanism of MC/DC analyzer, so this will not exercise the actual result / real values for MC/DC. This is a serious issue to fix, and we present a solution in this paper. We develop a dynamic version of MC/DC analyzer which actually takes care of the updated values of variables present in different predicates at different locations according to the appearance in a program. We will present an example for both Static and Dynamic analyses to measure MC/DC. Also, we compare the differences and handle those issues appearing in the analyses. We have observed that Dynamic mechanism has a significant advantage over static mechanism. We can achieve a significant improvement in MC/DC percentage, which is a big advantage of this proposed approach.

The rest of the article is organized as follows: Section 2 presents the Background Concepts. Section 3 presents the proposed approach Dy-COPECA. Section 4 shows the result analysis using an example. Section 5 compares our approach with existing work. Section 6 concludes the paper and suggests some future work.

2 BACKGROUND CONCEPTS

MC/DC is a code coverage criterion and was introduced by the RTCA DO-178B standard (Johnson et al., 1998). Test coverage approaches such as branch coverage which are popular for traditional programs are considered as inadequate for safety-critical systems. Thus, MC/DC was used to overcome this limitation and to achieve a linear growth of test case generation (Ammann et al., 2003; Bokil et al., 2009). MC/DC indicates that the outcome of a decision in the case of a conditional statement must be affected by the changes made to the individual conditions. MC/DC must satisfy the following criteria (Hayhurst, 2001):

- Every entry points and exit points of a program should be invoked at least once.
- Every decisions of a program should be invoked at least once for both true and false branch values.
- Every atomic conditions present in a decision should be invoked at least once for both true and false branches.
- Every possible outcomes of a decision must be affected by the changes made to each condition.

For example, let's take an example predicate "if(X && Y) then ...". In order to find the MC/DC test cases for this example, the following steps required to be performed:

Table 1: Extended truth table.

TC No.	X	Y	result	X	Y
TC ₁	True	True	True	TC ₃	TC ₂
TC ₂	True	False	False		TC ₁
TC ₃	False	True	False	TC ₁	
TC ₄	False	False	False		

- Create truth table for the predicate.
- Extend the truth table so that it indicates which test cases can be used to show the independence of each condition. The extended truth table is shown in Table 1.
- Show the independence of X by taking test cases {TC₁, TC₃} and the independence of Y by taking test case {TC₁, TC₂}.
- Union of the above test cases are known as the MC/DC test cases. The resulting MC/DC test cases are {TC₁, TC₂, and TC₃}, i.e., {(True, True)+(True, False)+(False, True)}

Definition 1. *Static MC/DC Analysis:* "Analyzing a C program to measure MC/DC using unit test data without any context consideration before the predicates."

Definition 2. *Dynamic MC/DC Analysis:* "Analyzing a C program to measure MC/DC using unit test data with considering the context before the predicate, which helps to process all the updated values of variables present in the whole program."

Definition 3. *Modified Condition / Decision Coverage:* "MC/DC is code coverage technique, where Condition is a Leaf level Boolean expression and Decision controls the program flow. MC/DC% is defined as the total number of independently affected conditions (I) out of total conditions (C) present in a program (Hayhurst, 2001; Hayhurst and Veerhusen, 2001)."

$$MCDC\% = \frac{|I|}{|C|} * 100\% \quad (1)$$

We have developed a tool called Dy-COPECA for measuring MC/DC%. Dy-COPECA stands for Dynamic Coverage PERcentage CALculator that accepts generated unit test cases along with C program and produces MC/DC%.

3 PROPOSED APPROACH: Dy-COPECA

In this section, we discuss the detailed description of Dy-COPECA using schematic representation.

Fig. 1 explains the schematic representation of Dy-COPECA. Dy-COPECA produces MC/DC% as

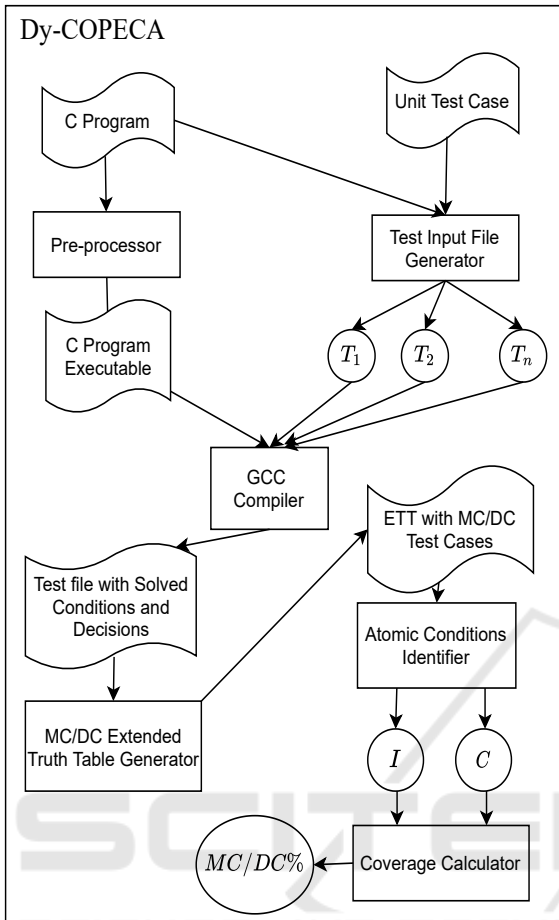


Figure 1: Schematic representation of Dy-COPECA.

output after imparting a C-program and unit test data as inputs. Dy-COPECA consists of six modules. These are (1) Pre-processor, (2) Test Input File Generator, (3) GCC Compiler, (4) MC/DC Extended Truth Table Generator, (5) Atomic Conditions Identifier, and (6) Coverage Calculator.

Now, let us discuss the functionalities and flow of execution for these modules. The flow starts with supplying C-program into *Pre-processor* as input. *Pre-processor* creates an annotated C-program and supplied into *GCC Compiler* to produce an executable code with functionalities processing test data files (*.a.out*). On the other hand, we have unit test data which generated from any of the test data generator (Younis et al., 2008) such as symbolic tester or concolic tester for C program. The test data generation is beyond the scope of this paper, so we assumed that we have already generated unit test data.

In practical, a unit test data file has several features for variables such as, variable's name, variable's size, and variable's input value etc. So, *Test Input File Generator* processes C-program along with unit test

Table 2: Assumed test data for example program in Fig. 2.

Var	T ₁	T ₂	T ₃	T ₄	T ₅	T ₆	T ₇	T ₈	T ₉	T ₁₀
p	21	21	15	15	80	85	60	70	30	30
q	45	51	45	51	40	35	55	60	55	85
r	0	90	110	90	71	60	25	69	50	80
s	0	0	10	0	50	125	90	119	80	80

data as inputs which separates all test input values in different files to get executed individually one by one.

Now, the third Module *GCC Compiler*, which is the actual reason for dynamic nature of our MC/DC analyzer. So, from Fig. 1, we can observe that, annotated C program code is imparted along with all test input files one by one into *GCC compiler* as inputs which provides an executable program which produces a text file with all solved condition and decisions after ran naively. *GCC compiler* executes C program at run time, which allots all concrete input values corresponding to all variables present in the program. It solves all the atomic conditions, decisions, arithmetic operations, assignments etc. But, our main focus is to solve atomic conditions and decisions for MC/DC mechanism.

Next, *MC/DC Extended Truth Table Generator*¹ module reads a text file as an input and produces ETT Table with MC/DC Test Cases as an output as shown in Fig. 1 according to the definition explained in Section 2.

Now, *Atomic Conditions Identifier* module reads ETT Table with MC/DC Test Cases to identify the set of Independently affected atomic conditions (I) along with total number of atomic conditions (C). The last module *Coverage Calculator* takes the value of I and C as inputs, and uses the formula given in Eq. 1 to measure MC/DC%. Finally Dy-COPECA produces MC/DC% for the input C-program.

4 EXPERIMENTAL RESULT ANALYSIS

In this section we discuss about static and dynamic analyses of MC/DC using an example.

Let us take an example of C program as shown in Fig. 2. Also, assume that unit test data is available with us to compute MC/DC as shown in Table 2. First of all, this program has five variables, out of which four variables are generated automatically from test case generator and one variable is initialized in program itself. This program has two predicates, in which first predicate has three atomic conditions, second predicate has four atomic conditions.

¹It is to be noted that computation of ETT table does not follow short-circuiting properties.

Table 3: Extended Truth Table (ETT) for first predicate of example C program in Fig. 2 using Static and Dynamic Analyses.

Test Cases	(p>20)	(q<50)	(r>100)	$P_1=((p>20)\&\&((q<50)\ \ (r>100)))$	(p>20)	(q<50)	(r>100)
T ₁	TRUE	TRUE	FALSE	TRUE		T ₂ ,T ₇ ,T ₈ ,T ₉ ,T ₁₀	
T ₂	TRUE	FALSE	FALSE	FALSE		T ₁ , T ₅ ,T ₆	
T ₃	FALSE	TRUE	TRUE	FALSE			
T ₄	FALSE	FALSE	FALSE	FALSE			
T ₅	TRUE	TRUE	FALSE	TRUE		T ₂ ,T ₇ ,T ₈ ,T ₉ ,T ₁₀	
T ₆	TRUE	TRUE	FALSE	TRUE		T ₂ ,T ₇ ,T ₈ ,T ₉ ,T ₁₀	
T ₇	TRUE	FALSE	FALSE	FALSE		T ₁ , T ₅ ,T ₆	
T ₈	TRUE	FALSE	FALSE	FALSE		T ₁ , T ₅ ,T ₆	
T ₉	TRUE	FALSE	FALSE	FALSE		T ₁ , T ₅ ,T ₆	
T ₁₀	TRUE	FALSE	FALSE	FALSE		T ₁ , T ₅ ,T ₆	

```

1. #include<stdio.h>
2. int main(){
3.   int p,q,r,s,x=10;
4.   if((p>20)&&((q<50)\|(r>100))){
5.     x=x+50;
6.     printf("This is if-branch of 1st
7.     predicate");}
8.   else {
9.     x=x+70;
10.    printf("This is else-branch of 1st
11.    predicate");}
12.   if(((p<=x)&&(q<x))\|\|(r>70)&&(s<120))){
13.     printf("This is if-branch of 2nd
14.     predicate");}
15.   else{
16.     printf("This is else-branch of 2nd
17.     predicate");}
18.   return 0;
19. }
```

Figure 2: An example C program.

4.1 Static MC/DC Analysis

We know that static MC/DC analysis executes test cases and measures MC/DC% for predicates one by one. With this fact, let us now create Extended Truth Table (ETT) for both predicates. We start with predicate “ $P_1=((p>20)\&\&((q<50)\|\|(r>100)))$ ”. Here, we have three atomic conditions $\{(p>20), (q<50), \text{ and } (r>100)\}$. Using test data from Table 2, one by one we create ETT as shown in Table 3². Here, $(q<50)$ is the only condition which independently affects the outcome of whole predicate after toggling it value. The minimum test cases must be “n+1”, where n=1 in this case. There are two MC/DC test cases out of eight test cases for this first predicate.

Now we create ETT for the second predicate i.e.,

²Please note that, Table 3 is common for both Static MC/DC analysis and Dynamic MC/DC Analysis, which shows the ETT for first predicate. So, we have not drawn two separate Tables for both analyses.

“ $P_2=((p<=x)\&\&(q<x))\|\|(r>70)\&\&(s<120))$ ”.

Since, this is a static analysis in which, it only able to process the test input values and available in test data generated. For P_2 , we have four atomic conditions $\{(p<=x), (q<x), (r>70), \text{ and } (s<120)\}$. Invoking this predicate from C-program, we must have test input values for p,q,r , and s variables, according to test data generated. We do not have the updated value of “x” variable, due to which we can not execute $(p<=x)$ and $(q<x)$. But, now it is possible to compute the outcome of $(r>70)$ and $(s<120)$ atomic conditions. It doesn’t mean that we are able to measure MC/DC%. According to problem, the run-time value $(p<=x)$ and $(q<x)$ cannot be computed and these are highly connected to $(r>70)$ and $(s<120)$ through a Boolean operator “ $\|\|$ ”. Hence the whole predicate is not able to compute MC/DC. Therefore, the MC/DC percentage for C program as shown in Fig. 2. using static MC/DC analysis can be computed as follows:

- Total number of atomic conditions present in the example C-program (C) = 7 $\{(p>20), (q<50), (r>100), (p<=x), (q<x), (r>70), \text{ and } (s<120)\}$.
- Total number of independently affected atomic conditions present in predicates of a C program (I) = 1 $\{(q<50)\}$.
- MC/DC test cases = $\{T_1 = (21,45,0,0)$ and $T_9 = (30,55,50,80)\}$. [Minimum number of test cases]³
- Using Eq. 1, we achieved 14.28% MC/DC.

4.2 Dynamic MC/DC Analysis

In this section, we discuss dynamic MC/DC analysis for the C-program shown Fig. 2. We know that

³We can choose either of the pair from the Table 3 as MC/DC test cases to show that $(q<50)$ is independently affected atomic condition.

dynamic MC/DC analyzer executes test cases and measure MC/DC% for predicates one by one by taking care of all the values of variables according to the procedural control flow of C-program. For more clarity we can observe the binary execution tree and all possible explored paths in Figures 3 and 4 respectively. We create Extended Truth Table (ETT) for both the predicates starting with “ $P_1=((p>20)\&\&((q<50)\| (r>100)))$ ”. Here we may observe that the ETT for “ P_1 ” is same as static MC/DC analysis which is already shown in Table 3. So, we are not going to discuss it again. But, in the second predicate “ $P_2(((p<=x)\&\&(q<x)\| ((r>70)\&\&(s<120))))$ ”, which has four atomic conditions $\{(p<=x), (q<x), (r>70), \text{ and } (s<120)\}$, we can observe that variable “ x ” is used in first two atomic conditions. Also, please observe Figures 2 and 3 where the variable “ x ” has different values at different appearance in the program. According to run-time execution, we have different values of the variable “ x ” at different appearance, and hence $(p<=x)$, and $(q<x)$ able to show the independently affected atomic conditions. Additionally, the atomic condition $(r>70)$ is also shown as independently affected atomic condition. The Extended Truth Table (ETT) for this predicate of example C-program using dynamic analysis is shown in Table 4. Therefore, the MC/DC percentage for C-program as shown in Fig. 2 using Dynamic MC/DC analysis can be computed as follows:

- Total number of atomic conditions present in predicates of the example C-program (C) = 7 $\{(p>20), (q<50), (r>100), (p<=x), (q<x), (r>70), \text{ and } (s<120)\}$.
- Total number of independently affected atomic conditions present the program (I) = 4 $\{(q<50), (p<=x), (q<x), \text{ and } (r>70)\}$.
- MC/DC test cases = $\{T_1=(21,45,0,0), T_4=(15,51,90,0), T_6=(85,35,60,125), T_9=(30,55,50,80) \text{ and } T_{10}=(30,85,80,80)\}$. [Minimum number of test cases show 4 atomic conditions (n+1)]
- Using Eq. 1, we computed 57.14% MC/DC.

4.3 Inference of Analysis

We have observed from our first analysis i.e., Static MC/DC analysis that the obtained MC/DC for C-program is 14.28% with two MC/DC test cases $\{T_1, \text{ and } T_9\}$. On the other hand, in our second analysis i.e. Dynamic MC/DC analysis, which processed all predicates and checked for all seven atomic conditions.

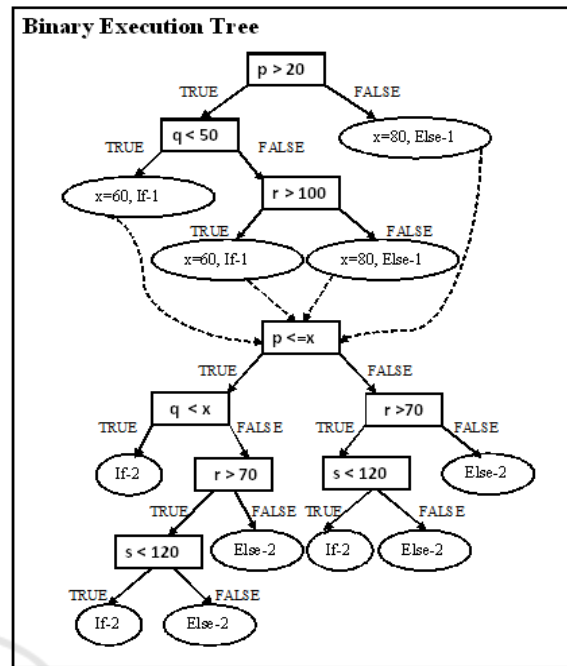


Figure 3: Binary Execution Tree in the example program given in Fig. 2.

1. FALSE-FALSE-FALSE
2. FALSE-FALSE-TRUE-FALSE
3. FALSE-FALSE-TRUE-TRUE
4. FALSE-TRUE-FALSE-FALSE
5. FALSE-TRUE-FALSE-TRUE-FALSE
6. FALSE-TRUE-FALSE-TRUE-TRUE
7. FALSE-TRUE-TRUE
8. TRUE-FALSE-FALSE-FALSE-FALSE
9. TRUE-FALSE-FALSE-FALSE-TRUE-FALSE
10. TRUE-FALSE-FALSE-FALSE-TRUE-TRUE
11. TRUE-FALSE-FALSE-TRUE-FALSE-FALSE
12. TRUE-FALSE-FALSE-TRUE-FALSE-TRUE-FALSE
13. TRUE-FALSE-FALSE-TRUE-FALSE-TRUE-TRUE
14. TRUE-FALSE-FALSE-TRUE-TRUE
15. TRUE-FALSE-TRUE-FALSE-FALSE
16. TRUE-FALSE-TRUE-FALSE-TRUE-FALSE
17. TRUE-FALSE-TRUE-FALSE-TRUE-TRUE
18. TRUE-FALSE-TRUE-TRUE-FALSE-FALSE
19. TRUE-FALSE-TRUE-TRUE-FALSE-TRUE-FALSE
20. TRUE-FALSE-TRUE-TRUE-FALSE-TRUE-TRUE
21. TRUE-FALSE-TRUE-TRUE-TRUE
22. TRUE-TRUE-FALSE-FALSE
23. TRUE-TRUE-FALSE-TRUE-FALSE
24. TRUE-TRUE-FALSE-TRUE-TRUE
25. TRUE-TRUE-TRUE-FALSE-FALSE
26. TRUE-TRUE-TRUE-FALSE-TRUE-FALSE
27. TRUE-TRUE-TRUE-FALSE-TRUE-TRUE
28. TRUE-TRUE-TRUE-TRUE

Figure 4: All possible paths explored in the example program given in Fig. 2.

Table 4: Extended Truth Table (ETT) for Second predicate of example C program in Fig. 2 using Dynamic analysis.

Test Cases	(p<=x)	(q<x)	(r>70)	(s<120)	$P_2 = (((p <= x) \& \& (q < x)) \& \& ((r > 70) \& \& (s > 120)))$	(p<=x)	(q<x)	(r>70)	(s<120)
T ₁	TRUE	TRUE	FALSE	TRUE	TRUE		T ₉	T ₁₀	
T ₂	TRUE	TRUE	TRUE	TRUE	TRUE				
T ₃	TRUE	TRUE	TRUE	TRUE	TRUE				
T ₄	TRUE	TRUE	TRUE	FALSE	TRUE		T ₆		
T ₅	FALSE	TRUE	TRUE	TRUE	TRUE				
T ₆	FALSE	TRUE	TRUE	FALSE	FALSE		T ₄		
T ₇	TRUE	TRUE	FALSE	TRUE	TRUE				
T ₈	TRUE	TRUE	FALSE	TRUE	TRUE				
T ₉	TRUE	FALSE	FALSE	TRUE	FALSE		T ₁		
T ₁₀	TRUE	FALSE	TRUE	TRUE	FALSE			T ₁	

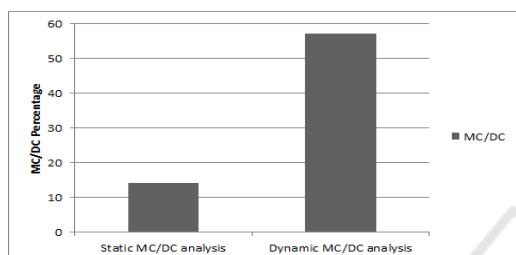


Figure 5: Comparison of analyses.

Dynamic analysis has improved the MC/DC percentage and reported 57.14% with four atomic conditions as Independently affecting conditions. For this, there are five test cases $\{T_1, T_4, T_6, T_9, \text{ and } T_{10}\}$ are required to compute MC/DC percentage of whole program. Due to Dynamic nature of MC/DC analysis, we achieve 42.88% of MC/DC higher as compared to static MC/DC analysis. The comparison between Static MC/DC analysis and Dynamic MC/DC analysis is shown in Fig. 5. We agree that, we have not achieved 100% MC/DC for C-program due to less number of test data assumed/available. Once we improve the test data, so we may achieve higher MC/DC.

4.4 Threats to Validity

1. Since, we focused on MC/DC percentage, programs without predicates are not useful for our experimental study.
2. In a predicate, there should be at least two conditions, because for MC/DC we require at least one logical operator.
3. We have not experimented a C-program which have multiple files to be invoked by main file, so current version may not handle such type of programs. We will implement this feature in future work.
4. Dy-COPECA is unable to solve the issues of redundant test data, in result it computes all the test

cases which are some times not required. Dy-COPECA required extra features to identify duplicate test data.

5 COMPARISON WITH RELATED WORK

In this section, we discuss some existing related work on this topic.

Hayhurst et al. (Hayhurst, 2001; Hayhurst and Veerhusen, 2001) presented a tutorial on MC/DC approach for aviation software applications that must comply with regulatory guided for Do-178B/C level A software. They have provided a five steps process to determine MC/DC, using that verification analyst recommend for the certification. From the tutorial available online, it shows that how to determine MC/DC for single predicate or decision. Also, this tutorial do not reflect any automation of the process. Dy-COPECA is actually fully automated and also able to execute a complete program which may consists of any number of predicates.

Chang et al. (Chang and Huang, 2007) proposed and developed a practical regression testing tool TASTE (Tool fo Automatic Software regression TESting). In their paper, they have used an useful method which focuses on all conditions of Boolean expression to determine MC/DC. Their proposed approach used *n-cube graph* and *gray code* to implement the MC/DC criterion. They have differentiated the necessary and redundant test cases. Also, their approach is dynamic in nature. Like, Chang et al., we also proposed Dy-COPECA which measure MC/DC. Dy-COPECA is dynamic nature to compute MC/DC at run-time, as like TASTE. Chang et al. have proposed some strategies to generate test cases, whereas our approach is more generic and can be plugged to any test cases generator.

Kandl et al. (Kandl and Kirner, 2010) implemented MC/DC for automotive domain. They have targeted

to inspect the error-detection rate of a set of test that attain higher possible MC/DC coverage. They initiated by generation of test cases followed error detection. Dy-COPECA is not targeted for error detection according to MC/DC. Also, we assumed for the unit test cases.

Ghani et al.(Ghani and Clark, 2009) introduce a search based testing technique to generate test case and using those test case they have computed MC/DC automatically. Their tool has advantage to compute Multiple Condition Coverage (MCC) and MC/DC. They have used *simulated annealing optimization* technique. It is not very clear from the paper that Ghani et al. focuses on dynamic nature or not. In our proposed approach, we have not used any test case generator technique, but can be extended in future work. Also, Dy-COPECA can not determine MCC. But, it uses dynamic analysis to compute MC/DC.

Like Ghani et al.(Ghani and Clark, 2009), Awedekian et al.(Awedekian et al., 2009) implemented search based algorithm for MC/DC test cases generation. Awedekian et al.(Awedekian et al., 2009) adopted the *Hill Climbing(HC)* and *Genetic Algorithm (GA)* to implement their approach. On the other hand, Dy-COPECA do not use any test cases generated technique. Only it follows the definition of standard *Unique cause MC/DC* provided by Hayhurst et al.(Hayhurst, 2001) at NASA. We applied dynamic behavior on the definition to automate the coverage tool.

Haque et al.(Haque et al., 2014) have proposed and developed a tool called MC/DC GEN. They have considered the issue of “Masking” in MC/DC. Based on masking MC/DC, they have implemented tool to generate test cases and determine MC/DC. But, their tool does not explain the dynamic nature, only they execute single predicate. In our proposed work, we are focusing on “Unique cause MC/DC” only. Masking MC/DC is beyond the scope of this paper. But, the major advantage of Dy-COPECA is its dynamic nature over Haque et al.(Haque et al., 2014) work.

Godbole et al.(Godbole et al., 2018b; Godbole et al., 2018a; Godbole et al., 2017a; Godbole et al., 2017b; Godbole et al., 2016; Godbole et al., 2013b; Godbole et al., 2013a; Godbole, 2013) have proposed several techniques to compute MC/DC using concolic test case generation technique. Their MC/DC analyzers were computing MC/DC score statically i.e. processing one predicate at a time. On the other hand, Dy-COPECA focused for the dynamic behavior of the program for MC/DC.

Table 5 shows the comparison with related work. All the related works considered focused on computing MC/D criterion. Our main objective of this paper

is to show the dynamic behavior is good to achieve higher MC/DC. So, we have taken the comparison factor as the mechanism of computing MC/DC either static or dynamic. We can observe that only Chang et al. have worked other than us on dynamic MC/DC analysis, and other have taken static MC/DC analysis. Based on our investigation from the paper, we presented that static is a very serious issue which needed to be fixed. Hence, we have proposed and developed a tool to overcome this issue.

Table 5: Comparison with related work.

Author's Name	Static MC/DC	Dynamic MC/DC
Hayhurst et al.(Hayhurst, 2001)	✓	X
Chang et al.(Chang and Huang, 2007)	X	✓
Kandl et al.(Kandl and Kirner, 2010)	✓	X
Ghani et al.(Ghani and Clark, 2009)	✓	X
Awedekian et al.(Awedekian et al., 2009)	✓	X
Haque et al.(Haque et al., 2014)	✓	X
Our Proposed work: Dy-COPECA	X	✓

6 CONCLUSION AND FUTURE WORK

We have proposed and developed Dy-COPECA, which is used to measure MC/DC at run time when we supply a C-program along with unit test cases. We have shown the limitations of static nature of coverage tool. Due to this issue, MC/DC gets degraded for a C-program. Also, we have explained the Dynamic analysis to improve the results. Using an example C-program the result analysis is explained. We have improved 42.88% of MC/DC through dynamic MC/DC analysis as compared to static MC/DC analysis.

In future work, we try to plug this Dy-COPECA with some test case generator tools such as symbolic tester, and concolic tester. Also, we plan to implement a Java version of Dy-COPECA.

REFERENCES

- Ammann, P., Offutt, J., and Huang, H. (2003). Coverage criteria for logical expressions. In *14th International Symposium on Software Reliability Engineering, 2003. ISSRE 2003.*, pages 99–107. IEEE.
- Awedekian, Z., Ayari, K., and Antoniol, G. (2009). Mc/dc automatic test input data generation. In *Proceedings of the 11th Annual conference on Genetic and evolutionary computation*, pages 1657–1664.
- Beizer, B. (2003). *Software testing techniques*. Dreamtech Press.
- Bird, D. L. and Munoz, C. U. (1983). Automatic generation of random self-checking test cases. *IBM systems journal*, 22(3):229–245.

- Bokil, P., Darke, P., Shrotri, U., and Venkatesh, R. (2009). Automatic test data generation for c programs. In *2009 Third IEEE International Conference on Secure Software Integration and Reliability Improvement*, pages 359–368. IEEE.
- Chang, J.-R. and Huang, C.-Y. (2007). A study of enhanced mc/dc coverage criterion for software testing. In *31st Annual International Computer Software and Applications Conference (COMPSAC 2007)*, volume 1, pages 457–464. IEEE.
- Chauhan, N. (2010). *Software Testing: Principles and Practices*. Oxford university press.
- Chilenski, J. J. and Miller, S. P. (1994). Applicability of modified condition/decision coverage to software testing. *Software Engineering Journal*, 9(5):193–200.
- Csallner, C., Smaragdakis, Y., and Xie, T. (2008). Dsd-crasher: A hybrid analysis tool for bug finding. *ACM Transactions on Software Engineering and Methodology (TOSEM)*, 17(2):1–37.
- DeMillo, R. A. and Offutt, A. J. (1993). Experimental results from an automatic test case generator. *ACM Transactions on Software Engineering and Methodology (TOSEM)*, 2(2):109–127.
- Gao, J., Espinoza, R., and He, J. (2005). Testing coverage analysis for software component validation. In *29th Annual International Computer Software and Applications Conference (COMPSAC'05)*, volume 1, pages 463–470. IEEE.
- Ghani, K. and Clark, J. A. (2009). Automatic test data generation for multiple condition and mc/dc coverage. In *2009 Fourth International Conference on Software Engineering Advances*, pages 152–157. IEEE.
- Godbole, S. (2013). *Improved modified condition/decision coverage using code transformation techniques*. PhD thesis.
- Godbole, S., Dutta, A., Mohapatra, D. P., Das, A., and Mall, R. (2016). Making a concolic tester achieve increased mc/dc. *Innovations in systems and software engineering*, 12(4):319–332.
- Godbole, S., Dutta, A., Mohapatra, D. P., and Mall, R. (2017a). J3 model: a novel framework for improved modified condition/decision coverage analysis. *Computer Standards & Interfaces*, 50:1–17.
- Godbole, S., Dutta, A., Mohapatra, D. P., and Mall, R. (2018a). Gecojap: A novel source-code preprocessing technique to improve code coverage. *Computer Standards & Interfaces*, 55:27–46.
- Godbole, S., Dutta, A., Mohapatra, D. P., and Mall, R. (2018b). Scaling modified condition/decision coverage using distributed concolic testing for java programs. *Computer Standards & Interfaces*, 59:61–86.
- Godbole, S., Mohapatra, D. P., Das, A., and Mall, R. (2017b). An improved distributed concolic testing approach. *Software: Practice and Experience*, 47(2):311–342.
- Godbole, S., Prashanth, G., Mohapatra, D. P., and Majhi, B. (2013a). Enhanced modified condition/decision coverage using exclusive-nor code transformer. In *2013 International Multi-Conference on Automation, Computing, Communication, Control and Compressed Sensing (iMac4s)*, pages 524–531. IEEE.
- Godbole, S., Prashanth, G., Mohapatra, D. P., and Majhi, B. (2013b). Increase in modified condition/decision coverage using program code transformer. In *2013 3rd IEEE International Advance Computing Conference (IACC)*, pages 1400–1407. IEEE.
- Grindal, M., Offutt, J., and Andler, S. F. (2005). Combination testing strategies: a survey. *Software Testing, Verification and Reliability*, 15(3):167–199.
- Gupta, N., Mathur, A. P., and Soffa, M. L. (1998). Automated test data generation using an iterative relaxation method. *ACM SIGSOFT Software Engineering Notes*, 23(6):231–244.
- Haque, A., Khalil, I., and Zamli, K. Z. (2014). An automated tool for mc/dc test data generation. In *2014 IEEE Symp. Comput. Informatics, Kota Kinabalu, Sabah, Malaysia*.
- Hayhurst, K. J. (2001). *A practical tutorial on modified condition/decision coverage*. DIANE Publishing.
- Hayhurst, K. J. and Veerhusen, D. S. (2001). A practical approach to modified condition/decision coverage. In *20th DASC. 20th Digital Avionics Systems Conference (Cat. No. 01CH37219)*, volume 1, pages 1B2–1. IEEE.
- Johnson, L. A. et al. (1998). Do-178b, software considerations in airborne systems and equipment certification. *Crosstalk, October*, 199:11–20.
- Jones, J. A. and Harrold, M. J. (2003). Test-suite reduction and prioritization for modified condition/decision coverage. *IEEE Transactions on software Engineering*, 29(3):195–209.
- Kandl, S. and Kirner, R. (2010). Error detection rate of mc/dc for a case study from the automotive domain. In *IFIP International Workshop on Software Technologies for Embedded and Ubiquitous Systems*, pages 131–142. Springer.
- Majumdar, R. and Sen, K. (2007). Hybrid concolic testing. In *29th International Conference on Software Engineering (ICSE'07)*, pages 416–426. IEEE.
- Mall, R. (2018). *Fundamentals of software engineering*. PHI Learning Pvt. Ltd.
- Myers, G. J., Sandler, C., and Badgett, T. (2011). *The art of software testing*. John Wiley & Sons.
- Rajan, A., Whalen, M. W., and Heimdahl, M. P. (2008). The effect of program and model structure on mc/dc test adequacy coverage. In *Proceedings of the 30th international conference on Software engineering*, pages 161–170.
- Younis, M. I., Zamli, K. Z., and Isa, N. A. M. (2008). Irps—an efficient test data generation strategy for pairwise testing. In *International Conference on Knowledge-Based and Intelligent Information and Engineering Systems*, pages 493–500. Springer.